

the witness reduction technique

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presentation structure

- motivation
- preliminaries
- what *The witness reduction technique* is about
- warm up
- a central question

presentation structure

- other questions...

- about **OptP** and **SpanP**

- resources

motivation

*This technique is mainly used in order to **relate potential closure**
properties of #P to complexity class hierarchy collapses!*

preliminaries: classes

language classes

□ NP, *co*NP

□ PH

□ UP

□ \oplus P

function classes

□ FP

□ #P

□ OptP

□ SpanP

preliminaries: classes

language classes (continued)

□ SPP

□ PP

□ C=P

preliminaries: operations

□ **addition** $a + b$

□ **multiplication** $a \cdot b$

□ **proper subtraction** $a \ominus b = \max(\{0, a - b\})$

□ **proper division** $a \oslash b = \lfloor a/b \rfloor$

preliminaries: operations

- **proper decrement** $a \ominus 1$
- **proper division by two** $a \oslash 2$
- **maximum** $\max(\{a, b\}) = a$ if $a \geq b$ else b
- **minimum** $\min(\{a, b\}) = a$ if $a \leq b$ else b

(pause)

$\text{Closed}(\blacksquare, \boxplus) \equiv$ the complexity class \blacksquare is closed under the operation \boxplus

(pause)

Turing Machine \equiv **Non-Deterministic** and **Polynomial-Time** Turing
Machine

warm up

Closed($\#P$, +)

$f, g \in \#P$

$$\exists M_1, M_2 : f(x) = \#acc_{M_1}(x) \wedge g(x) = \#acc_{M_2}(x)$$

$$M_3(x) =_{\text{def}} M_1(x) \text{ or } M_2(x)$$

$$\#acc_{M_3}(x) = f(x) + g(x)$$

$$\#acc_{M_3}(x) \in \#P \Rightarrow f(x) + g(x) \in \#P$$

warm up

Closed($\#P, \cdot$)

$f, g \in \#P$

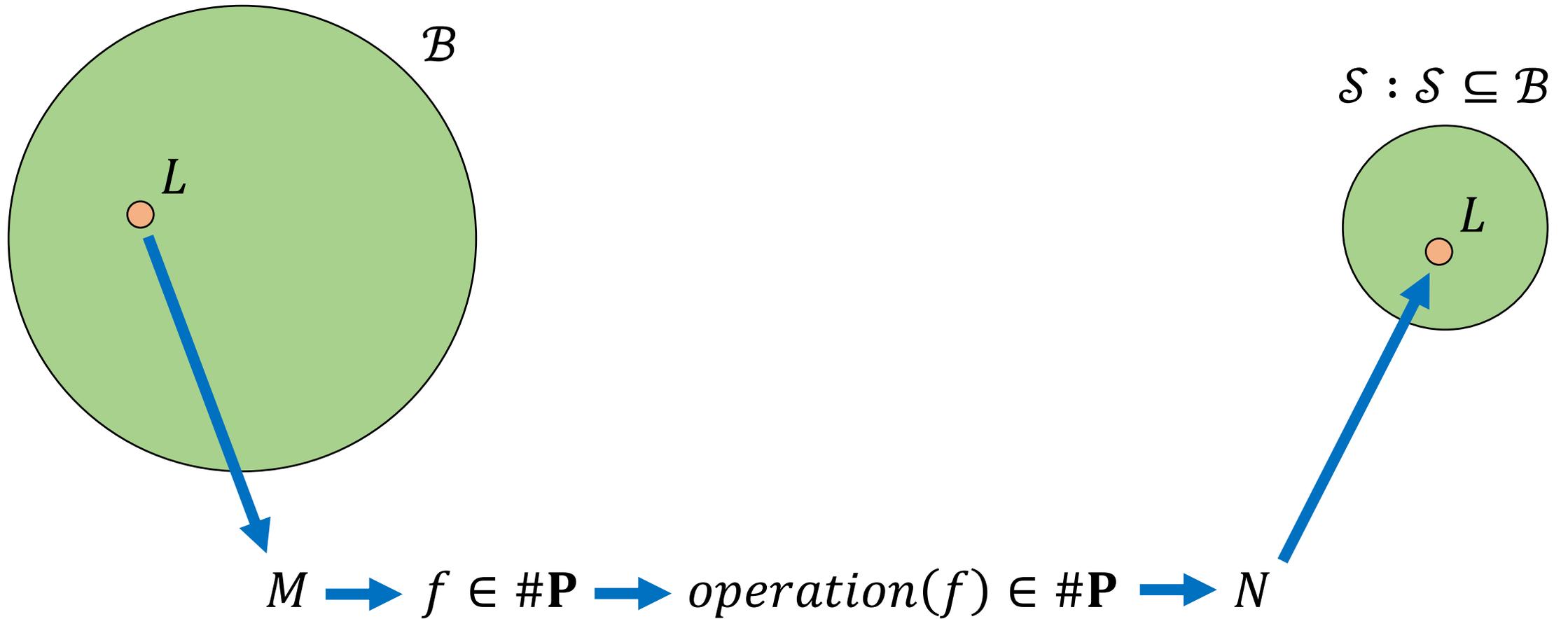
$$\exists M_1, M_2 : f(x) = \#acc_{M_1}(x) \wedge g(x) = \#acc_{M_2}(x)$$

$$M_3(x) =_{\text{def}} M_1(x) \text{ and } M_2(x)$$

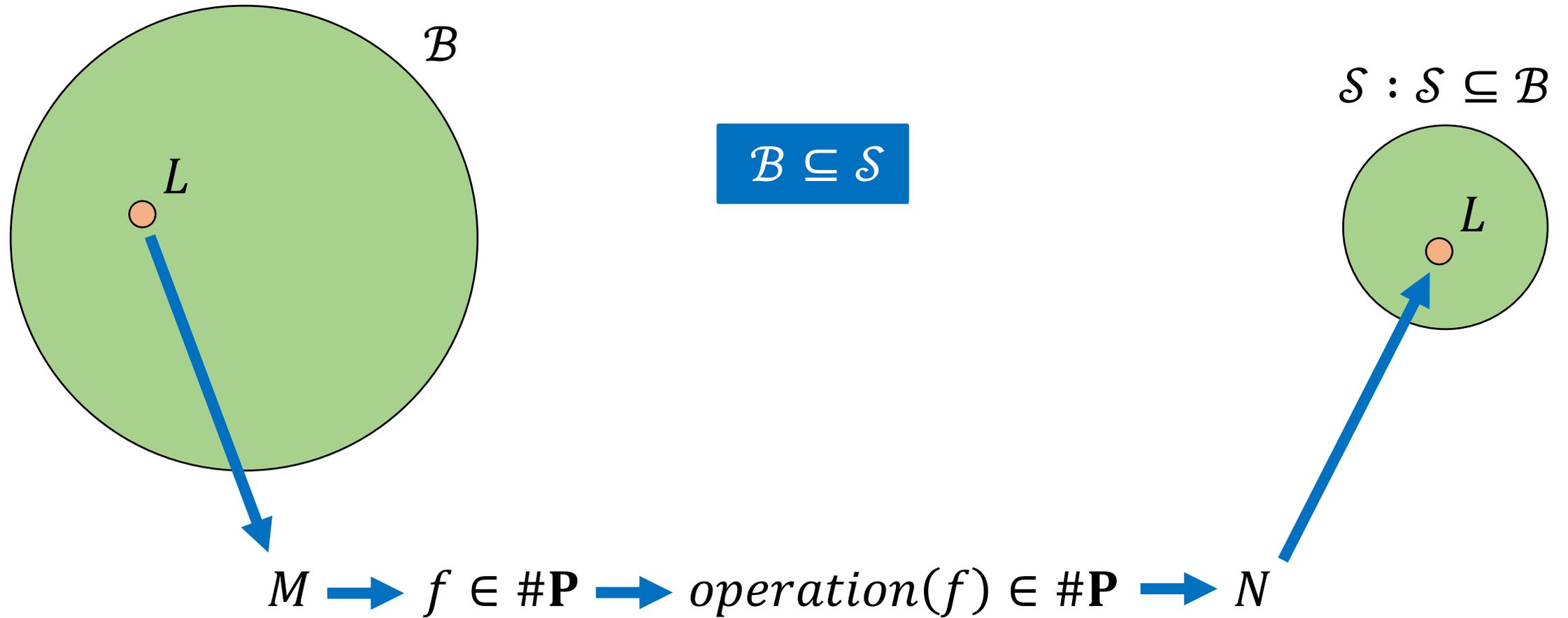
$$\#acc_{M_3}(x) = f(x) \cdot g(x)$$

$$\#acc_{M_3}(x) \in \#P \Rightarrow f(x) \cdot g(x) \in \#P$$

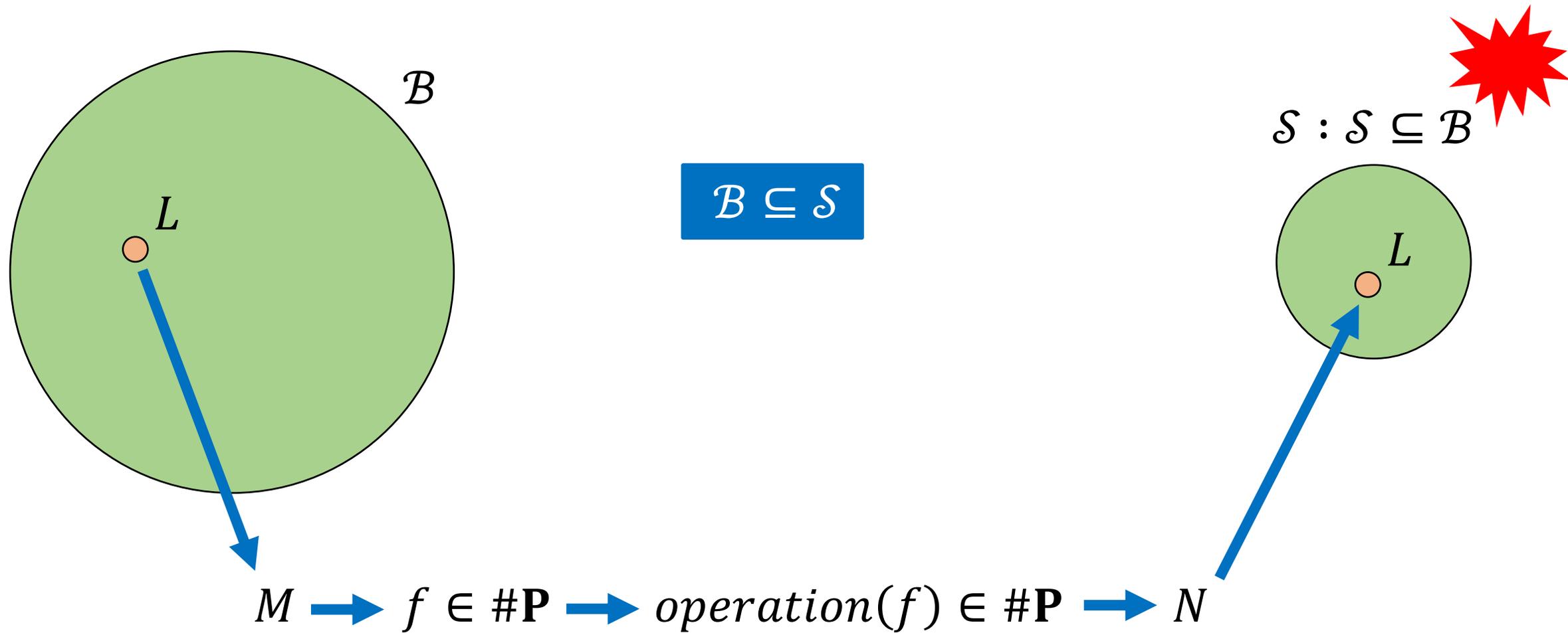
what *The witness reduction technique* is about



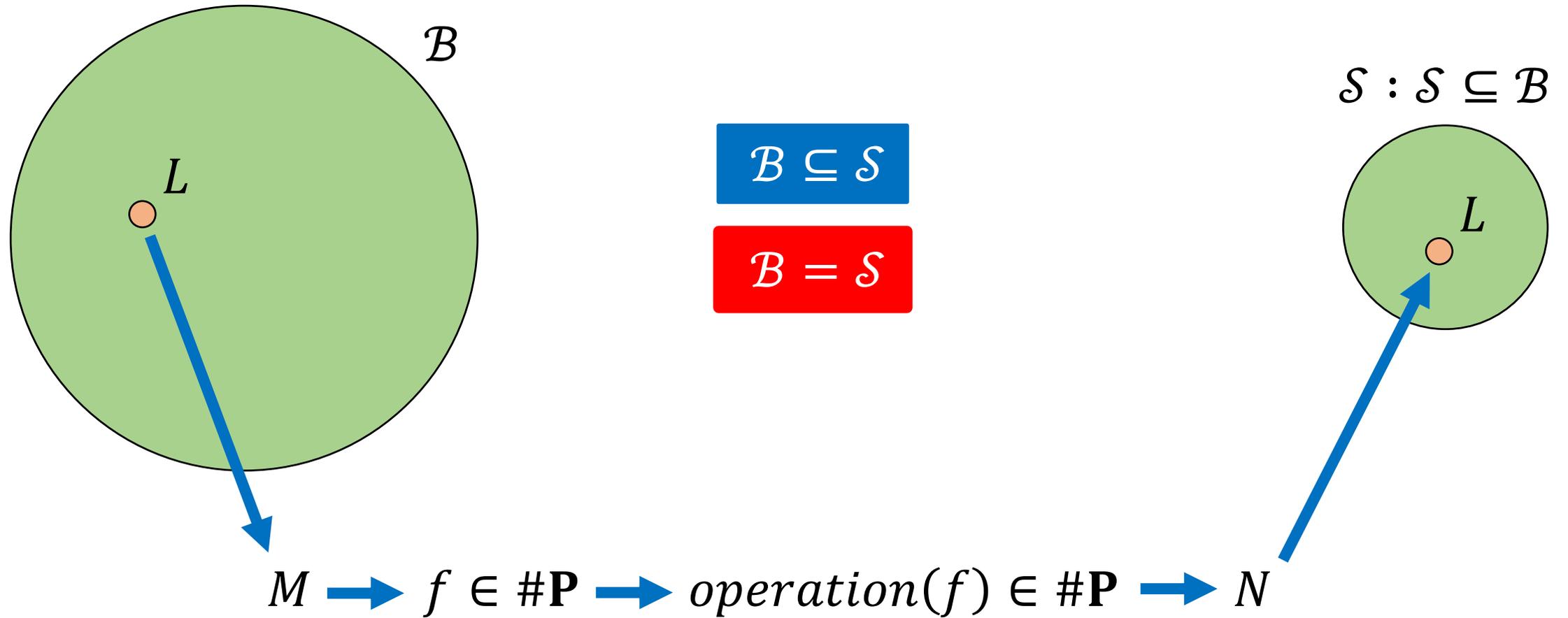
what *The witness reduction technique* is about



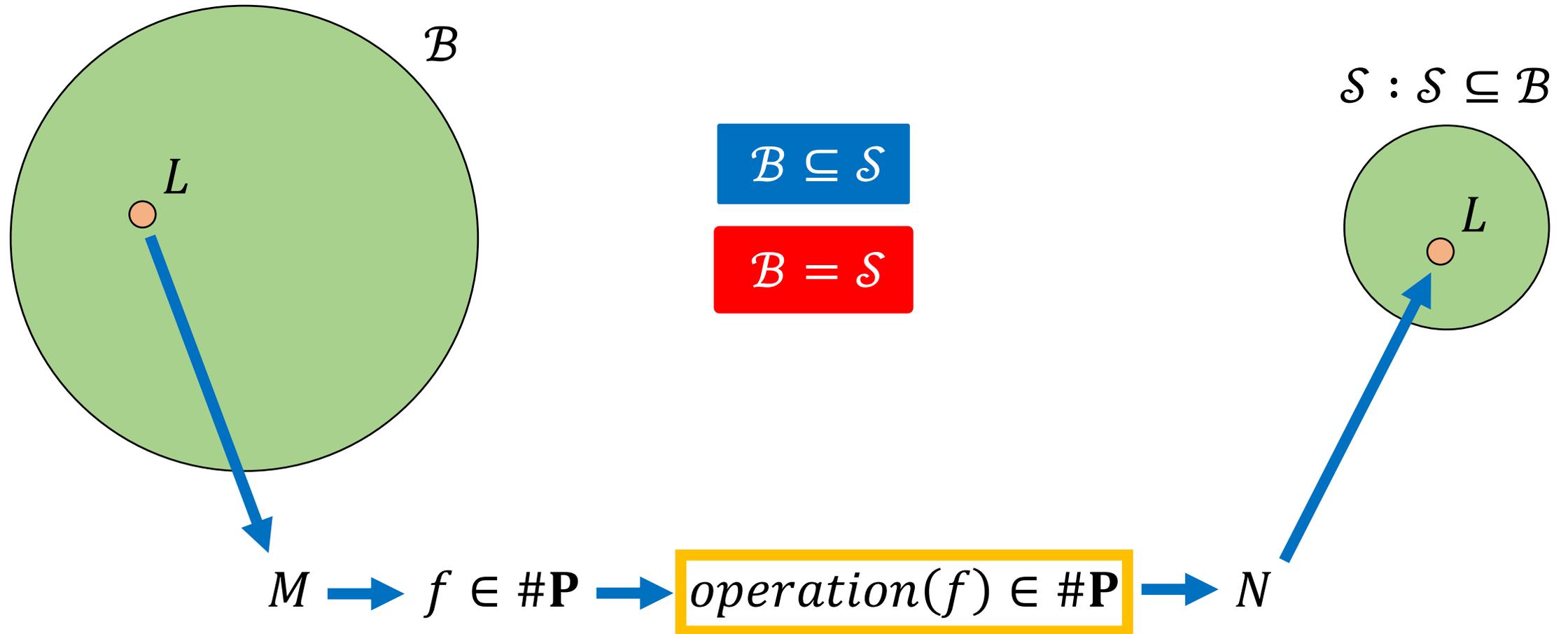
what *The witness reduction technique* is about



what *The witness reduction technique* is about



what *The witness reduction technique* is about



a central question

Closed($\#P, \Theta$) =? **True**

a central question

1. $\text{Closed}(\#\mathbf{P}, \Theta)$
2. $\text{Closed}(\#\mathbf{P}, op)$
3. **UP = PP**

$op : \text{Function}(op, \mathbb{N}^2, \mathbb{N}) \wedge \text{PolyTime}(op)$

$$1 \equiv 2 \equiv 3$$

a central question: a theorem

1. $\text{Closed}(\#\mathbf{P}, \Theta)$
2. $\text{Closed}(\#\mathbf{P}, op)$
3. $\mathbf{UP} = \mathbf{PP}$

$op : \text{Function}(op, \mathbb{N}^2, \mathbb{N}) \wedge \text{PolyTime}(op)$

$$1 \equiv 2 \equiv 3$$

Proof

$$1 \rightarrow 3 \quad 3 \rightarrow 2 \quad 2 \rightarrow 1$$

(pause)

FP \subseteq #P

a theorem: proof, 1 \rightarrow 3

Closed($\#P, \Theta$) \rightarrow **UP = PP**

a theorem: proof, 1 \rightarrow 3

Closed($\#P, \Theta$) \rightarrow **UP = PP**

Closed($\#P, \Theta$) \rightarrow **PP \subseteq UP \wedge UP \subseteq PP**

UP \subseteq PP : True

Closed($\#P, \Theta$) \rightarrow **PP \subseteq UP**

Closed($\#P, \Theta$) \rightarrow **PP \subseteq coNP \wedge coNP \subseteq UP**

a theorem: proof, 1 \rightarrow 3

Closed($\#P, \Theta$) \rightarrow $PP \subseteq coNP \wedge coNP \subseteq UP$

$\Rightarrow \begin{cases} \text{Closed}(\#P, \Theta) \rightarrow PP \subseteq coNP \\ \text{Closed}(\#P, \Theta) \rightarrow coNP \subseteq UP \end{cases}$

a theorem: proof, 1 \rightarrow 3

$$\begin{cases} \text{Closed}(\#P, \Theta) \rightarrow \mathbf{PP} \subseteq \mathbf{coNP} \\ \text{Closed}(\#P, \Theta) \rightarrow \mathbf{coNP} \subseteq \mathbf{UP} \end{cases}$$

a theorem: proof, 1 \rightarrow 3

Closed($\#P, \Theta$) \rightarrow $PP \subseteq \text{coNP}$

$L \in PP$

$$\Rightarrow L = \left\{ x : |\{y : R(x, y) \wedge |y| \leq p(|x|)\}| \geq \frac{2^{p(|x|)}}{2} + 1 \right\}$$

a theorem: proof, 1 \rightarrow 3

Turing Machine $M(x)$

guess $y : |y| \leq p(|x|)$

return 1 if $R(x, y)$ else 0

a theorem: proof, 1 \rightarrow 3

Turing Machine $M(x)$

$$\Rightarrow L(M) = L$$

$$\& \exists f : f(x) = \#acc_M(x) \in \#\mathbf{P}$$

guess $y : |y| \leq p(|x|)$

return 1 if $R(x, y)$ **else** 0

$$\text{observe} \begin{cases} x \in L \rightarrow f(x) \geq \frac{2^{p(|x|)}}{2} + 1 \\ x \notin L \rightarrow f(x) \leq \frac{2^{p(|x|)}}{2} \end{cases}$$

a theorem: proof, 1 \rightarrow 3

$$g(x) = \frac{2^{p(|x|)}}{2} \in \mathbf{FP} \wedge \mathbf{FP} \subseteq \#\mathbf{P} \Rightarrow g(x) \in \#\mathbf{P}$$

$$\text{Closed}(\#\mathbf{P}, \Theta) \Rightarrow h(x) = f(x) \Theta g(x) \in \#\mathbf{P}$$

$$\Rightarrow \exists N : h(x) = \#acc_N(x)$$

a theorem: proof, 1 \rightarrow 3

$$x \in L \rightarrow f(x) \geq \frac{2^{p(|x|)}}{2} + 1 \ \& \ g(x) = \frac{2^{p(|x|)}}{2}$$

$$\rightarrow h(x) = f(x) \ominus g(x) \geq 1$$

$$x \notin L \rightarrow f(x) \leq \frac{2^{p(|x|)}}{2} \ \& \ g(x) = \frac{2^{p(|x|)}}{2}$$

$$\rightarrow h(x) = f(x) \ominus g(x) = 0$$

a theorem: proof, 1 \rightarrow 3

$$x \in L \rightarrow h(x) \geq 1$$

$$x \notin L \rightarrow h(x) = 0$$

a theorem: proof, 1 \rightarrow 3

$$\text{so } \begin{cases} x \in L \rightarrow h(x) \geq 1 \\ x \notin L \rightarrow h(x) = 0 \end{cases} \text{ and } \exists N : h(x) = \#acc_N(x)$$

$\Rightarrow L(N) = L \ \& \ L \in \mathbf{NP}$

$\Rightarrow \mathbf{PP} \subseteq \mathbf{NP}$

$\Rightarrow \mathbf{coPP} \subseteq \mathbf{coNP}$

$\mathbf{coPP} = \mathbf{PP} \Rightarrow \mathbf{PP} \subseteq \mathbf{coNP}$

a theorem: proof, 1 \rightarrow 3

$\left\{ \begin{array}{l} \text{Closed}(\#P, \Theta) \rightarrow \mathbf{PP} \subseteq \mathbf{coNP} \text{ OK} \\ \text{Closed}(\#P, \Theta) \rightarrow \mathbf{coNP} \subseteq \mathbf{UP} \end{array} \right.$

a theorem: proof, 1 \rightarrow 3

Closed($\#\mathbf{P}, \Theta$) \rightarrow $\text{coNP} \subseteq \mathbf{UP}$

$L \in \text{coNP}$

$\Rightarrow \bar{L} \in \mathbf{NP}$

$\Rightarrow \exists M : L(M) = \bar{L}$

$\Rightarrow \exists f : f(x) = \#acc_M(x) \in \#\mathbf{P}$

a theorem: proof, $1 \rightarrow 3$

so $\exists f : f(x) = \#acc_M(x) \in \#P$ and $L(M) = \bar{L}$

observe $\begin{cases} x \in \bar{L} \rightarrow f(x) \geq 1 \\ x \notin \bar{L} \rightarrow f(x) = 0 \end{cases} \Rightarrow \begin{cases} x \in L \rightarrow f(x) = 0 \\ x \notin L \rightarrow f(x) \geq 1 \end{cases}$

$1 \in FP \wedge FP \subseteq \#P \Rightarrow 1 \in \#P$

$\text{Closed}(\#P, \ominus) \Rightarrow g(x) = 1 \ominus f(x) \in \#P$

a theorem: proof, 1 \rightarrow 3

$$g(x) = 1 \ominus f(x) \in \#P$$

$$\Rightarrow \exists N : \#acc_N(x) = g(x)$$

$$\text{observe } \begin{cases} x \in L \rightarrow f(x) = 0 \rightarrow g(x) = 1 \ominus 0 = 1 \\ x \notin L \rightarrow f(x) \geq 1 \rightarrow g(x) = 1 \ominus f(x) = 0 \end{cases}$$

a theorem: proof, 1 \rightarrow 3

$$g(x) = 1 \ominus f(x) \in \#P$$

$$\Rightarrow \exists N : \#acc_N(x) = g(x)$$

$$\text{observe} \begin{cases} x \in L \rightarrow g(x) = 1 \\ x \notin L \rightarrow g(x) = 0 \end{cases}$$

$$\Rightarrow L(N) = L \ \& \ L \in \mathbf{UP}$$

$$\Rightarrow \mathbf{coNP} \subseteq \mathbf{UP}$$

a theorem: proof, 1 \rightarrow 3

$\left\{ \begin{array}{l} \text{Closed}(\#P, \Theta) \rightarrow \mathbf{PP} \subseteq \mathbf{coNP} \text{ OK} \\ \text{Closed}(\#P, \Theta) \rightarrow \mathbf{coNP} \subseteq \mathbf{UP} \text{ OK} \end{array} \right.$

a theorem: proof, 1 \rightarrow 3

$$\begin{cases} \text{Closed}(\#P, \Theta) \rightarrow \mathbf{PP} \subseteq \mathbf{coNP} \text{ OK} \\ \text{Closed}(\#P, \Theta) \rightarrow \mathbf{coNP} \subseteq \mathbf{UP} \text{ OK} \end{cases}$$

$\Rightarrow \mathbf{PP} \subseteq \mathbf{UP}$

$\mathbf{UP} \subseteq \mathbf{PP} \Rightarrow \mathbf{PP} = \mathbf{UP}$

a central question: a theorem

1. $\text{Closed}(\#\mathbf{P}, \Theta)$
2. $\text{Closed}(\#\mathbf{P}, op)$
3. **UP = PP**

$$1 \equiv 2 \equiv 3$$

Proof

$$1 \rightarrow 3 \quad 3 \rightarrow 2 \quad 2 \rightarrow 1$$

a theorem: proof, 3 \rightarrow 2

PP = UP \rightarrow Closed(**#P**, *op*)

op : Function(*op*, \mathbb{N}^2 , \mathbb{N}) \wedge PolyTime(*op*)

a theorem: proof, 3 \rightarrow 2

PP = UP \rightarrow Closed($\#P$, op)

$f, g \in \#P$

$$B_f = \{\langle x, n \rangle : f(x) \geq n\}$$

$$B_g = \{\langle x, n \rangle : g(x) \geq n\}$$

a theorem: proof, 3 \rightarrow 2

$$B_f = \{\langle x, n \rangle : f(x) \geq n\}$$

$$B_g = \{\langle x, n \rangle : g(x) \geq n\}$$

$B_f \in \text{PP}$ & $B_g \in \text{PP}$

a theorem: proof, 3 \rightarrow 2

V

$$= \{ \langle x, n_1, n_2 \rangle : \langle x, n_1 \rangle \in B_f \wedge \langle x, n_1 + 1 \rangle \notin B_f \wedge \langle x, n_2 \rangle$$

a theorem: proof, 3 \rightarrow 2

$B_f \in \mathbf{PP} \ \& \ B_g \in \mathbf{PP} \Rightarrow \text{DisjointUnion}(B_f, B_g) = B_f \uplus B_g \in \mathbf{PP}$

$$X \uplus Y = \{0x : x \in X\} \cup \{1y : y \in Y\}$$

a theorem: proof, 3 \rightarrow 2

$B_f \in \mathbf{PP} \ \& \ B_g \in \mathbf{PP} \Rightarrow \text{DisjointUnion}(B_f, B_g) = B_f \uplus B_g \in \mathbf{PP}$

$V \leq_{\text{bounded truth table}} B_f \uplus B_g$

a theorem: proof, 3 \rightarrow 2

$B_f \in \mathbf{PP} \ \& \ B_g \in \mathbf{PP} \Rightarrow \text{DisjointUnion}(B_f, B_g) = B_f \uplus B_g \in \mathbf{PP}$

$V \leq_{\text{bounded truth table}} B_f \uplus B_g$

$\text{Closed}(\mathbf{PP}, \leq_{\text{bounded truth table}}) \Rightarrow V \in \mathbf{PP}$

a theorem: proof, 3 \rightarrow 2

$$V \in \text{PP} \ \& \ \text{PP} = \text{UP} \Rightarrow V \in \text{UP}$$

a theorem: proof, 3 \rightarrow 2

$$V \in \text{PP} \ \& \ \text{PP} = \text{UP} \Rightarrow V \in \text{UP}$$

$$\Rightarrow \exists N : L(N) = V$$

$$\& \forall x, n_1, n_2 : \#acc_N(\langle x, n_1, n_2 \rangle) = 0 \text{ or } 1$$

a theorem: proof, 3 \rightarrow 2

V

$$= \{ \langle x, n_1, n_2 \rangle : \langle x, n_1 \rangle \in B_f \wedge \langle x, n_1 + 1 \rangle \notin B_f \wedge \langle x, n_2 \rangle$$

a theorem: proof, $3 \rightarrow 2$

Turing Machine $M(x)$

guess i, j

if $N(\langle x, i, j \rangle) = \text{yes}$ **then**

 guess $k : k \in \{1, \dots, op(i, j)\}$

 accept

return k

else reject

a theorem: proof, 3 \rightarrow 2

Turing Machine $M(x)$

guess i, j

if $N(\langle x, i, j \rangle) = \text{yes}$ then

 guess $k : k \in \{1, \dots, \text{op}(i, j)\}$

 accept

return k

else reject

$$\forall x \exists i, j : i = f(x) \ \& \ j = g(x)$$

$$\Rightarrow \forall x \exists i, j : \langle x, i, j \rangle \in V$$

$$\Rightarrow \forall x \exists i, j : N(\langle x, i, j \rangle) = \text{yes}$$

$$\Rightarrow \forall x \exists i, j : \#acc_N(\langle x, i, j \rangle) \geq 1$$

$$\Rightarrow \forall x \exists i, j : \#acc_N(\langle x, i, j \rangle) = 1$$

a theorem: proof, 3 \rightarrow 2

$$x, i, j : \#acc_N(\langle x, i, j \rangle) = 1$$

$$\Rightarrow N(\langle x, i, j \rangle) = \text{yes}$$

$$\Rightarrow \langle x, i, j \rangle \in V$$

$$\Rightarrow i = f(x) \wedge j = g(x)$$

$$\Rightarrow op(i, j) = op(f(x), g(x))$$

$$\Rightarrow k \in \{1, \dots, op(f(x), g(x))\}$$

a theorem: proof, 3 \rightarrow 2

Turing Machine $M(x)$

guess i, j

if $N(\langle x, i, j \rangle) = \text{yes}$ then

guess $k : k \in \{1, \dots, op(f(x), g(x))\}$

accept

return k

else reject

a theorem: proof, 3 \rightarrow 2

$$\#acc_M(x) = ?$$

a theorem: proof, 3 \rightarrow 2

$$\#acc_M(x) = \#acc_N(\langle x, i, j \rangle) \cdot op(f(x), g(x))$$

a theorem: proof, 3 \rightarrow 2

$$\#acc_M(x) = \#acc_N(\langle x, i, j \rangle) \cdot op(f(x), g(x))$$

$$\Rightarrow \#acc_M(x) = 1 \cdot op(f(x), g(x))$$

a theorem: proof, 3 \rightarrow 2

Turing Machine $M(x)$

guess i, j

if $N(\langle x, i, j \rangle) = \text{yes}$ **then**

 guess $k : k \in \{1, \dots, op(i, j)\}$

 accept

return k

else reject

$$\forall x \exists i, j : i = f(x) \ \& \ j = g(x)$$

$$\Rightarrow \forall x \exists i, j : \langle x, i, j \rangle \in V$$

$$\Rightarrow \forall x \exists i, j : N(\langle x, i, j \rangle) = \text{yes}$$

$$\Rightarrow \forall x \exists i, j : \#acc_N(\langle x, i, j \rangle) \geq 1$$

$$\Rightarrow \forall x \exists i, j : \#acc_N(\langle x, i, j \rangle) = 1$$

a theorem: proof, 3 \rightarrow 2

$$\#acc_M(x) = \#acc_N(\langle x, i, j \rangle) \cdot op(f(x), g(x))$$

$$\Rightarrow \#acc_M(x) = 1 \cdot op(f(x), g(x)) = op(f(x), g(x))$$

$$\Rightarrow op(f(x), g(x)) \in \#P$$

$$\Rightarrow \text{Closed}(\#P, op)$$

a central question: a theorem

1. $\text{Closed}(\#P, \Theta)$
2. $\text{Closed}(\#P, op)$
3. **UP = PP**

$$1 \equiv 2 \equiv 3$$

Proof

$$1 \rightarrow 3 \quad 3 \rightarrow 2 \quad 2 \rightarrow 1$$

a theorem: proof, 2 \rightarrow 1

$\text{Closed}(\#\mathbf{P}, op) \rightarrow \text{Closed}(\#\mathbf{P}, \Theta)$

$op : \text{Function}(op, \mathbb{N}^2, \mathbb{N}) \wedge \text{PolyTime}(op)$

a theorem: proof, 2 \rightarrow 1

$\text{Closed}(\#\mathbf{P}, op) \rightarrow \text{Closed}(\#\mathbf{P}, \Theta)$

True

$op : \text{Function}(op, \mathbb{N}^2, \mathbb{N}) \wedge \text{PolyTime}(op)$

a central question: a theorem

1. $\text{Closed}(\#\mathbf{P}, \Theta)$
2. $\text{Closed}(\#\mathbf{P}, op)$
3. **UP = PP**

$$1 \equiv 2 \equiv 3$$

Proof

$$1 \rightarrow 3 \quad 3 \rightarrow 2 \quad 2 \rightarrow 1$$

a central question: a theorem

1. $\text{Closed}(\#P, \Theta)$
2. $\text{Closed}(\#P, op)$
3. **UP = PP**

$$1 \equiv 2 \equiv 3$$

Proved!

a central question: consequences

UP = PP \rightarrow

UP = NP = coNP = PH = \bigoplus P = PP \cup PP^{PP} \cup PP^{PP^{PP}} \cup ...

an alternative way

1. Closed(**#P**, \bigcirc)
2. Closed(**#P**, op)
3. **UP = PP**

$$op : \text{Function}(op, \mathbb{N}^2, \mathbb{N}) \wedge \text{PolyTime}(op)$$

$$1 \equiv 2 \equiv 3$$

we can use \bigcirc instead of \ominus

other questions...

What are the potential closure properties of $\#\mathbf{P}$ that do not imply that

$\#\mathbf{P}$ is closed under any polynomial-time operation?

other questions...: some implications

$\text{Closed}(\#P, \ominus 1) \rightarrow \text{coNP} \subseteq \text{SPP}$

$\text{Closed}(\#P, \oslash 2) \rightarrow \bigoplus P = \text{SPP}$

$\text{Closed}(\#P, \text{min}) \rightarrow \text{NP} = \text{UP}$

$\text{Closed}(\#P, \text{max}) \vee \text{Closed}(\#P, \text{min}) \rightarrow \text{C}_=P = \text{SPP}$

Closed($\#P, \ominus 1$) \rightarrow $\text{coNP} \subseteq \text{SPP}$

$L \in \text{NP}$

$\Rightarrow \exists M : L(M) = L$

$\Rightarrow \exists f : f(x) = \#acc_M(x) \in \#P$

$1 \in \text{FP} \wedge \text{FP} \subseteq \#P \Rightarrow 1 \in \#P$

Closed($\#P, \ominus$) $\Rightarrow g(x) = f(x) \ominus 1 \in \#P$

$\Rightarrow \exists N : g(x) = \#acc_N(x)$

Closed($\#P, \ominus 1$) \rightarrow $\text{coNP} \subseteq \text{SPP}$

$$\exists N : g(x) = f(x) \ominus 1 = \#acc_N(x)$$

$$N'(x) = \neg N(x)$$

$$\Rightarrow \#acc_{N'}(x) = \#rej_N(x) = 2^{p(|x|)} - g(x) \in \#P$$

$$\text{Closed}(\#P, +) \Rightarrow h(x) = f(x) + \left(2^{p(|x|)} - g(x)\right) \in \#P$$

$$\Rightarrow \exists K : \#acc_K(x) = h(x) = f(x) + 2^{p(|x|)} - (f(x) \ominus 1)$$

Closed($\#P, \ominus 1$) \rightarrow $\text{coNP} \subseteq \text{SPP}$

$$\exists K : \#acc_K(x) = h(x) = f(x) + 2^{p(|x|)} - (f(x) \ominus 1)$$

$$x \notin L \Rightarrow f(x) = 0 \Rightarrow f(x) \ominus 1 = 0 \Rightarrow h(x) = 2^{p(|x|)}$$

$$x \in L \Rightarrow f(x) > 0 \Rightarrow f(x) \ominus 1 = f(x) - 1 \Rightarrow h(x) = 2^{p(|x|)} + 1$$

Closed($\#P, \ominus 1$) \rightarrow $\text{coNP} \subseteq \text{SPP}$

$$\exists K : \#acc_K(x) = h(x) = f(x) + 2^{p(|x|)} - (f(x) \ominus 1)$$

$$x \notin L \Rightarrow h(x) = 2^{p(|x|)}$$

$$x \in L \Rightarrow h(x) = 2^{p(|x|)} + 1$$

$$\Rightarrow L(K) = L \ \& \ L \in \mathbf{SPP}$$

$$\Rightarrow \mathbf{NP} \subseteq \mathbf{SPP}$$

$$\Rightarrow \mathbf{coNP} \subseteq \mathbf{coSPP}$$

Closed($\#P, \Theta 1$) \rightarrow $coNP \subseteq SPP$

so $coNP \subseteq coSPP$

$coSPP = SPP \Rightarrow coNP \subseteq SPP$

other questions...: some implications

$\text{Closed}(\#P, \ominus 1) \rightarrow \text{coNP} \subseteq \text{SPP}$ **OK**

$\text{Closed}(\#P, \oslash 2) \rightarrow \oplus P = \text{SPP}$

$\text{Closed}(\#P, \text{min}) \rightarrow \text{NP} = \text{UP}$

$\text{Closed}(\#P, \text{max}) \vee \text{Closed}(\#P, \text{min}) \rightarrow \text{C}_=P = \text{SPP}$

Closed($\#P, \emptyset, 2$) $\rightarrow \oplus P = SPP$

$\oplus P \subseteq SPP \wedge SPP \subseteq \oplus P$

$L \in \oplus P$

$\Rightarrow \exists M : L(M) = L$

$\Rightarrow \exists f : f(x) = \#acc_M(x) \in \#P$

$2 \in FP \wedge FP \subseteq \#P \Rightarrow 2 \in \#P$

Closed($\#P, \emptyset$) $\Rightarrow f(x) \emptyset 2 \in \#P$

Closed($\#P, \bigoplus 2$) $\rightarrow \bigoplus P = SPP$

$$\text{Closed}(\#P, \cdot) \Rightarrow 2 \cdot (f(x) \bigoplus 2) \in \#P$$

$$g(x) = \#rej_M(x) = 2^{p(|x|)} - \#acc_M(x) = 2^{p(|x|)} - f(x)$$

$$\Rightarrow g(x) = \#acc_{\neg M}(x) \in \#P$$

$$\text{Closed}(\#P, +) \Rightarrow h(x) = g(x) + 2 \cdot (f(x) \bigoplus 2) \in \#P$$

$$\Rightarrow \exists N : \#acc_N(x) = h(x) = 2^{p(|x|)} - f(x) + 2 \cdot (f(x) \bigoplus 2)$$

Closed($\#P, \emptyset 2$) $\rightarrow \oplus P = SPP$

$$h(x) = 2^{p(|x|)} - f(x) + 2 \cdot (f(x) \emptyset 2)$$

$$x \notin L \Rightarrow \text{Even}(f(x)) \Rightarrow 2 \cdot (f(x) \emptyset 2) = f(x) \Rightarrow h(x) = 2^{p(|x|)}$$

$$\begin{aligned} x \in L \Rightarrow \text{Odd}(f(x)) \Rightarrow 2 \cdot (f(x) \emptyset 2) &= f(x) - 1 \Rightarrow h(x) \\ &= 2^{p(|x|)} - 1 \end{aligned}$$

Closed($\#P, \emptyset, 2$) $\rightarrow \bigoplus P = SPP$

$$h(x) = 2^{p(|x|)} - f(x) + 2 \cdot (f(x) \emptyset 2)$$

$$x \notin L \rightarrow h(x) = 2^{p(|x|)}$$

$$x \in L \rightarrow h(x) = 2^{p(|x|)} - 1$$

$\Rightarrow L(N) = L \ \& \ L \in coSPP$

$\Rightarrow \bigoplus P \subseteq coSPP$

$coSPP = SPP \Rightarrow \bigoplus P \subseteq SPP$

Closed($\#P, \emptyset, 2$) $\rightarrow \oplus P = SPP$

$$\oplus P \subseteq SPP \wedge SPP \subseteq \oplus P$$

True

Closed($\#P, \emptyset, 2$) $\rightarrow \oplus P = SPP$

$$\left. \begin{array}{l} \oplus P \subseteq SPP \\ SPP \subseteq \oplus P \end{array} \right\} \Rightarrow \oplus P = SPP$$

other questions...: some implications

$\text{Closed}(\#P, \ominus 1) \rightarrow \text{coNP} \subseteq \text{SPP}$ **OK**

$\text{Closed}(\#P, \oslash 2) \rightarrow \bigoplus P = \text{SPP}$ **OK**

$\text{Closed}(\#P, \min) \rightarrow \text{NP} = \text{UP}$

$\text{Closed}(\#P, \max) \vee \text{Closed}(\#P, \min) \rightarrow \text{C}_=P = \text{SPP}$

Closed(#P, min) \rightarrow NP = UP

Closed(#P, min) \rightarrow NP \subseteq UP \wedge UP \subseteq NP

$L \in \text{NP}$

$\Rightarrow \exists M : L(M) = L$

$\Rightarrow \exists f : f(x) = \#acc_M(x) \in \#P$

$1 \in \text{FP} \wedge \text{FP} \subseteq \#P \Rightarrow 1 \in \#P$

Closed(#P, min) $\Rightarrow g(x) = \min(\{f(x), 1\}) \in \#P$

Closed(#P, min) \rightarrow NP = UP

$$g(x) = \min(\{f(x), 1\}) \in \#P$$

$$\Rightarrow \exists N : \#acc_N(x) = g(x)$$

$$x \in L \rightarrow f(x) \geq 1 \rightarrow g(x) = 1$$

$$x \notin L \rightarrow f(x) = 0 \rightarrow g(x) = 0$$

Closed($\#P$, min) \rightarrow NP = UP

$$g(x) = \min(\{f(x), 1\}) \in \#P$$

$$\Rightarrow \exists N : \#acc_N(x) = g(x)$$

$$x \in L \rightarrow g(x) = 1$$

$$x \notin L \rightarrow g(x) = 0$$

$$\Rightarrow L(N) = L \ \& \ L \in \mathbf{UP}$$

$$\Rightarrow \mathbf{NP} \subseteq \mathbf{UP}$$

Closed(#P, min) \rightarrow NP = UP

Closed(#P, min) \rightarrow NP \subseteq UP \wedge UP \subseteq NP

True

Closed(#P, min) \rightarrow NP = UP

$$\left. \begin{array}{l} \text{NP} \subseteq \text{UP} \\ \text{UP} \subseteq \text{NP} \end{array} \right\} \Rightarrow \text{NP} = \text{UP}$$

other questions...: some implications

Closed($\#P, \ominus 1$) \rightarrow $\text{coNP} \subseteq \text{SPP}$ **OK**

Closed($\#P, \oslash 2$) $\rightarrow \oplus P = \text{SPP}$ **OK**

Closed($\#P, \text{min}$) $\rightarrow \text{NP} = \text{UP}$ **OK**

Closed($\#P, \text{max}$) \vee Closed($\#P, \text{min}$) $\rightarrow \text{C}_=P = \text{SPP}$

other questions...: some implications

Closed($\#P, \ominus 1$) \rightarrow $\text{coNP} \subseteq \text{SPP}$ **OK**

Closed($\#P, \otimes 2$) \rightarrow $\oplus P = \text{SPP}$ **OK**

Closed($\#P, \text{min}$) \rightarrow $\text{NP} = \text{UP}$ **OK**

Closed($\#P, \text{max}$) \vee Closed($\#P, \text{min}$) \rightarrow $\text{C}_=P = \text{SPP}$

about OptP and SpanP

about OptP

1. Closed(**OptP**, Θ)
2. Closed(**OptP**, op)
3. **NP** = $coNP$

$op : \text{Function}(op, \mathbb{N}^2, \mathbb{N}) \wedge \text{PolyTime}(op)$

$$1 \equiv 2 \equiv 3$$

about **OptP**: theorem 2

1. Closed(**OptP**, Θ)
2. Closed(**OptP**, op)
3. **NP** = $coNP$

$op : \text{Function}(op, \mathbb{N}^2, \mathbb{N}) \wedge \text{PolyTime}(op)$

$$1 \equiv 2 \equiv 3$$

Proof

$$3 \rightarrow 2 \rightarrow 1 \rightarrow 3$$

(pause)

FP \subseteq OptP

theorem 2: proof, 3 \rightarrow 2

$\mathbf{NP} = \mathbf{coNP} \rightarrow \text{Closed}(\mathbf{OptP}, op)$

$op : \text{Function}(op, \mathbb{N}^2, \mathbb{N}) \wedge \text{PolyTime}(op)$

theorem 2: proof, 3 \rightarrow 2

$$\mathbf{NP} = \mathbf{coNP} \rightarrow \text{Closed}(\mathbf{OptP}, op)$$

$$f, g \in \mathbf{OptP}$$

$$\Rightarrow \exists N_f, N_g \forall x : f(x) = \text{MaxOutput}(N_f, x) \wedge$$

$$g(x) = \text{MaxOutput}(N_g, x)$$

theorem 2: proof, 3 \rightarrow 2

NP = coNP \rightarrow Closed(OptP, op)

$f, g \in \mathbf{OptP}$

$\Rightarrow \exists N_f, N_g \forall x : f(x) = \text{MaxOutput}(N_f, x) \wedge$

$g(x) = \text{MaxOutput}(N_g, x)$

$$L_f = \{\langle x, i \rangle : f(x) > i\}$$

$$L_g = \{\langle x, i \rangle : g(x) > i\}$$

theorem 2: proof, 3 \rightarrow 2

$\Rightarrow L_f, L_g \in \mathbf{NP}$

$\mathbf{NP} = \mathbf{coNP} \Rightarrow L_f, L_g \in \mathbf{coNP} \Rightarrow \overline{L_f}, \overline{L_g} \in \mathbf{NP}$

$$\overline{L_f} = \{\langle x, i \rangle : f(x) \leq i\}$$

$$\overline{L_g} = \{\langle x, i \rangle : g(x) \leq i\}$$

$\Rightarrow \exists N_1, N_2 : L(N_1) = \overline{L_f} \wedge L(N_2) = \overline{L_g}$

theorem 2: proof, 3 \rightarrow 2

Turing Machine $M(x)$

$$w_f = \text{AGuessedOutput}(N_f, x)$$

$$w_g = \text{AGuessedOutput}(N_g, x)$$

if $N_1(\langle x, w_f \rangle) \wedge N_2(\langle x, w_g \rangle)$ **then return** $op(w_f, w_g)$

else return 0

theorem 2: proof, 3 \rightarrow 2

Turing Machine $M(x)$

$$w_f = \text{AGuessedOutput}(N_f, x) \leq f(x) = \text{MaxOutput}(N_f, x)$$

$$w_g = \text{AGuessedOutput}(N_g, x) \leq g(x) = \text{MaxOutput}(N_g, x)$$

if $N_1(\langle x, w_f \rangle) \wedge N_2(\langle x, w_g \rangle)$ **then return** $op(w_f, w_g)$

else return 0

theorem 2: proof, 3 \rightarrow 2

$$N_1(\langle x, w_f \rangle) \wedge N_2(\langle x, w_g \rangle)$$

$$\Rightarrow f(x) \leq w_f \wedge g(x) \leq w_g \Rightarrow w_f \geq f(x) \wedge w_g \geq g(x)$$

$$\Rightarrow w_f = f(x) \wedge w_g = g(x)$$

$$\Rightarrow op(w_f, w_g) = op(f(x), g(x))$$

theorem 2: proof, 3 \rightarrow 2

Turing Machine $M(x)$

$w_f = \text{AnOutput}(N_f, x)$

$w_g = \text{AnOutput}(N_g, x)$

if $N_1(\langle x, w_f \rangle) \wedge N_2(\langle x, w_g \rangle)$ then return $op(f(x), g(x))$

else return 0

theorem 2: proof, 3 \rightarrow 2

$$\forall x \exists w_f, w_g : w_f = f(x) \ \& \ w_g = g(x)$$

$$\Rightarrow \forall x \exists w_f, w_g : N_1(\langle x, w_f \rangle) \wedge N_2(\langle x, w_g \rangle)$$

$$\Rightarrow \forall x \exists w_f, w_g : \text{the if condition is } \mathbf{True}$$

$$\Rightarrow \forall x \exists w_f, w_g : \text{Output}(M, x) = op(f(x), g(x))$$

theorem 2: proof, 3 \rightarrow 2

Turing Machine $M(x)$

$w_f = \text{AGuessedOutput}(N_f, x)$

$w_g = \text{AGuessedOutput}(N_g, x)$

if $N_1(\langle x, w_f \rangle) \wedge N_2(\langle x, w_g \rangle)$ then return $op(f(x), g(x))$

else return 0

theorem 2: proof, 3 \rightarrow 2

$$\Rightarrow \forall x : \text{MaxOutput}(M, x) = \text{op}(f(x), g(x))$$

$$\Rightarrow \text{op}(f(x), g(x)) \in \mathbf{OptP}$$

$$\Rightarrow \text{Closed}(\mathbf{OptP}, \text{op})$$

about **OptP**: theorem 2

1. Closed(**OptP**, Θ)
2. Closed(**OptP**, op)
3. **NP** = $coNP$

$$1 \equiv 2 \equiv 3$$

Proof

$$3 \rightarrow 2 \quad 2 \rightarrow 1 \quad 1 \rightarrow 3$$

theorem 2: proof, 2 \rightarrow 1

$\text{Closed}(\mathbf{OptP}, op) \rightarrow \text{Closed}(\mathbf{OptP}, \Theta)$

$op : \text{Function}(op, \mathbb{N}^2, \mathbb{N}) \wedge \text{PolyTime}(op)$

theorem 2: proof, 2 \rightarrow 1

$\text{Closed}(\mathbf{OptP}, op) \rightarrow \text{Closed}(\mathbf{OptP}, \Theta)$

True

$op : \text{Function}(op, \mathbb{N}^2, \mathbb{N}) \wedge \text{PolyTime}(op)$

about **OptP**: theorem 2

1. Closed(**OptP**, Θ)
2. Closed(**OptP**, op)
3. **NP** = $coNP$

$$1 \equiv 2 \equiv 3$$

Proof

$$3 \rightarrow 2 \quad 2 \rightarrow 1 \quad 1 \rightarrow 3$$

theorem 2: proof, 1 \rightarrow 3

$\text{Closed}(\mathbf{OptP}, \Theta) \rightarrow \mathbf{NP} = \mathbf{coNP}$

theorem 2: proof, 1 \rightarrow 3

Closed(OptP, Θ) \rightarrow NP = coNP

$L \in$ NP

$\Rightarrow \exists M : L(M) = L$

$\Rightarrow \text{MaxOutput}(M, x) = \begin{cases} 1, & \text{if } x \in L \\ 0, & \text{if } x \notin L \end{cases}$

$\Rightarrow \exists f : f(x) = \text{MaxOutput}(M, x) \in$ **OptP**

theorem 2: proof, 1 \rightarrow 3

$$1 \in \mathbf{FP} \wedge \mathbf{FP} \subseteq \mathbf{OptP} \Rightarrow 1 \in \mathbf{OptP}$$

$$\text{Closed}(\mathbf{OptP}, \ominus) \Rightarrow h(x) = 1 \ominus f(x) \in \mathbf{OptP}$$

$$\Rightarrow \exists N \forall x : \text{MaxOutput}(N, x) = h(x)$$

$$\text{observe } h(x) = \begin{cases} 1 \ominus 1 = 0, & \text{if } x \in L \\ 1 \ominus 0 = 1, & \text{if } x \notin L \end{cases}$$

theorem 2: proof, 1 \rightarrow 3

$$\text{so } h(x) = \begin{cases} 0, & \text{if } x \in L \\ 1, & \text{if } x \notin L \end{cases}$$

theorem 2: proof, 1 \rightarrow 3

$$\text{so } h(x) = \begin{cases} 0, & \text{if } x \in L \\ 1, & \text{if } x \notin L \end{cases} \Rightarrow h(x) = \begin{cases} 1, & \text{if } x \in \bar{L} \\ 0, & \text{if } x \notin \bar{L} \end{cases}$$

$$\Rightarrow L(N) = \bar{L} \text{ \& } \bar{L} \in \mathbf{NP}$$

$$\Rightarrow L \in \mathbf{coNP}$$

$$\Rightarrow \mathbf{NP} \subseteq \mathbf{coNP}$$

$$\Rightarrow \mathbf{coNP} \subseteq \mathbf{co(coNP)}$$

theorem 2: proof, 1 \rightarrow 3

so $\text{coNP} \subseteq \text{co}(\text{coNP})$

$\Rightarrow \text{coNP} \subseteq \text{NP}$

$$\left. \begin{array}{l} \text{NP} \subseteq \text{coNP} \\ \text{coNP} \subseteq \text{NP} \end{array} \right\} \Rightarrow \text{NP} = \text{coNP}$$

about **OptP**: theorem 2

1. Closed(**OptP**, Θ)
2. Closed(**OptP**, op)
3. **NP** = $coNP$

$$1 \equiv 2 \equiv 3$$

Proof

$$3 \rightarrow 2 \quad 2 \rightarrow 1 \quad 1 \rightarrow 3$$

about **OptP**: theorem 2

1. Closed(**OptP**, Θ)
2. Closed(**OptP**, op)
3. **NP** = co **NP**

$$1 \equiv 2 \equiv 3$$

Proved!

about SpanP

1. Closed(**SpanP**, Θ)
2. Closed(**SpanP**, op)
3. **NP = PH = PP^{NP}**

$op : \text{Function}(op, \mathbb{N}^2, \mathbb{N}) \wedge \text{PolyTime}(op)$

$$1 \equiv 2 \equiv 3$$

open issue

Find complexity classes \mathcal{A} and \mathcal{B} such that,

$$\text{Closed}(\#\mathbf{P}, \Theta 1) \leftrightarrow \mathcal{A} = \mathcal{B}$$

resources

- Hemaspaandra, L. A. and Ogihara, M., *The Complexity Theory Companion*, Springer, 2002, pp. 91-108

Thank(*You*)