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We want efficient algoriths for these Polynomial Operations (Multiplication, Addition, Evaluation)

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Some reperesentations $2 + 3x + x^2 \to A = [2, 3, 1]$ (-1,0),(0,2),(1,6)

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Fact: A polynomial of degree *d* can be uniquely represented by its values in any d+1 points. (at least d+1)

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Algorithms vs.	Representations		
	Coefficients	Roots	Samples
Evaluation	O(n)	O(n)	$O(n^2)$
Addition	O(n)	∞	O(n)
Multiplication	$O(n^2)$	O(n)	O(n)

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So we want $Coef \rightarrow Value$ representation Multiply and then $Value \rightarrow Coef$.

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NAINE APPROACH

$$\{(x_0, P(x_0)), (x_1, P(x_1)), \dots, (x_d, P(x_d))\}\}$$

$$P(x) = p_0 + p_1 x + p_2 x^2 + \dots + p_d x^d$$

$$P(x_0) = p_0 + p_1 x_0 + p_2 x_0^2 + \dots + p_d x_0^d$$

$$P(x_1) = p_0 + p_1 x_1 + p_2 x_1^2 + \dots + p_d x_1^d$$

$$\vdots$$

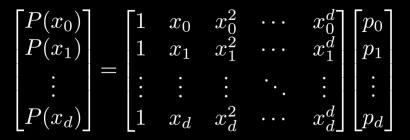
$$P(x_d) = p_0 + p_1 x_d + p_2 x_d^2 + \dots + p_d x_d^d$$

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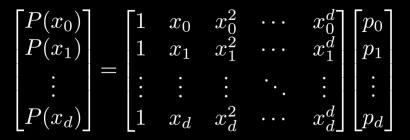
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To find the values in d+1 point computation is $O(d^2)$

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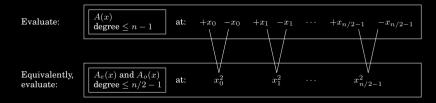
To find the values in d+1 point computation is $O(d^2)$

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What will happen when I compute $A(x_0)$ and $A(-x_0)$??? $3 + 4x + 6x^2 + 2x^3 + x^4 + 10x^5 = (3 + 6x^2 + x^4) + x(4 + 2x^2 + 10x^4).$ What will happen when I compute $P(x_0)$ and $P(-x_0)$??? $3 + 4x + 6x^2 + 2x^3 + x^4 + 10x^5 = (3 + 6x^2 + x^4) + x(4 + 2x^2 + 10x^4).$ $P(x_i) = P_e(x_i^2) + x_i P_o(x_i^2)$ $P(-x_i) = P_e(x_i^2) - x_i P_o(x_i^2)$

Evaluate only at squares of the original points (n/2).

The degree of P_e and P_o also drops to half!!



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$$\begin{split} & \begin{aligned} & \mathbb{E} \text{valuate} \begin{array}{c} \frac{P(x) : [p_0, p_1, \dots, p_{n-1}]}{[\pm x_1, \pm x_2, \dots, \pm x_{n/2}]} \\ & \mathbb{P}(x) = P_e(x^2) + xP_o(x^2) \end{aligned} \\ & \mathbb{E} \text{valuate} \begin{array}{c} \frac{P_e(x^2) : [p_0, p_2, \dots, p_{n-2}]}{[x_1^2, x_2^2, \dots, x_{n/2}^2]} \\ & \mathbb{E} \text{valuate} \begin{array}{c} \frac{P_o(x^2) : [p_1, p_3, \dots, p_{n-1}]}{[x_1^2, x_2^2, \dots, x_{n/2}^2]} \\ & \mathbb{E} \text{valuate} \begin{array}{c} \frac{P_o(x^2) : [p_1, p_3, \dots, p_{n-1}]}{[x_1^2, x_2^2, \dots, x_{n/2}^2]} \\ & \mathbb{E} \text{valuate} \begin{array}{c} P_o(x^2) : [p_1, p_3, \dots, p_{n-1}] \\ & [x_1^2, x_2^2, \dots, x_{n/2}^2] \end{array} \\ & \mathbb{E} \text{valuate} \begin{array}{c} P_o(x^2) : [p_1, p_3, \dots, p_{n-1}] \\ & [x_1^2, x_2^2, \dots, x_{n/2}^2] \end{array} \\ & \mathbb{E} \text{valuate} \begin{array}{c} P_o(x^2) : [p_1, p_3, \dots, p_{n-1}] \\ & [x_1^2, x_2^2, \dots, x_{n/2}^2] \end{array} \\ & \mathbb{E} \text{valuate} \begin{array}{c} P_o(x^2) : [p_1, p_3, \dots, p_{n-1}] \\ & [x_1^2, x_2^2, \dots, x_{n/2}^2] \end{array} \\ & \mathbb{E} \text{valuate} \begin{array}{c} P_o(x^2) : [p_0(x^2), P_0(x^2), \dots, P_o(x^2_{n/2})] \\ & \mathbb{E} \text{valuate} \begin{array}{c} P_o(x^2) : [p_0(x^2), P_0(x^2), \dots, P_o(x^2_{n/2})] \\ & \mathbb{E} \text{valuate} \begin{array}{c} P(x_1) = P_e(x_1^2) + x_1 P_0(x_1^2) \\ & P(-x_1) = P_e(x_1^2) - x_1 P_0(x_1^2) \\ & \mathbb{E} \text{valuate} \begin{array}{c} P(x_1) = P_e(x_1^2) + x_1 P_0(x_1^2) \\ & \mathbb{E} \text{valuate} \begin{array}{c} P(x_1) = P_e(x_1^2) + x_1 P_0(x_1^2) \\ & \mathbb{E} \text{valuat} \begin{array}{c} P(x_1) = P_e(x_1^2) + x_1 P_0(x_1^2) \\ & \mathbb{E} \text{valuat} \begin{array}{c} P(x_1) = P_e(x_1^2) + x_1 P_0(x_1^2) \\ & \mathbb{E} \text{valuat} \begin{array}{c} P(x_1) = P_e(x_1^2) + x_1 P_0(x_1^2) \\ & \mathbb{E} \text{valuat} \begin{array}{c} P(x_1) = P_e(x_1^2) + x_1 P_0(x_1^2) \\ & \mathbb{E} \text{valuat} \begin{array}{c} P(x_1) = P_e(x_1^2) + x_1 P_0(x_1^2) \\ & \mathbb{E} \text{valuat} \end{array} \end{aligned}$$

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It would be of a good complexity if there was no problem

$$T(n) = 2 \cdot T\left(\frac{n}{2}\right) + O(n) = O(n \lg n).$$

X the set of all points stops geting halfed after step 1.

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Solution: nth Roots of unity $z^n = 1$

Example 4th roots: $\{1, -1, i, -i\}$ Example 2nd roots: $\{1, -1\}$ Example 1st root: {1}

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Expand the domain of the polynomials to \mathbb{C} and everything still holds. With the help of some equations we denote the nth- roots by the complex numbers $1, \omega, \omega^2, \cdots, \omega^{n-1}$ where $\omega = e^{i2\pi/n}$. And pick $n = 2^l$. Expand the domain of the polynomials to \mathbb{C} and everything still holds. With the help of some equations we denote the nth- roots by the complex numbers $1, \omega, \omega^2, \cdots, \omega^{n-1}$ where $\omega = e^{i2\pi/n}$. And pick $n = 2^l$ for convinience.

$$\begin{bmatrix} P(x_0) \\ P(x_1) \\ P(x_2) \\ \vdots \\ P(x_{n-1}) \end{bmatrix} = \begin{bmatrix} 1 & x_0 & x_0^2 & \cdots & x_0^{n-1} \\ 1 & x_1 & x_1^2 & \cdots & x_1^{n-1} \\ 1 & x_2 & x_2^2 & \cdots & x_2^{n-1} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & x_{n-1} & x_{n-1}^2 & \cdots & x_{n-1}^{n-1} \end{bmatrix} \begin{bmatrix} p_0 \\ p_1 \\ p_2 \\ \vdots \\ p_{n-1} \end{bmatrix}$$

$$\begin{bmatrix} P(\omega^{0}) \\ P(\omega^{1}) \\ P(\omega^{2}) \\ \vdots \\ P(\omega^{n-1}) \end{bmatrix} = \begin{bmatrix} 1 & 1 & 1 & \cdots & 1 \\ 1 & \omega & \omega^{2} & \cdots & \omega^{n-1} \\ 1 & \omega^{2} & \omega^{4} & \cdots & \omega^{2(n-1)} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & \omega^{n-1} & \omega^{2(n-1)} & \cdots & \omega^{(n-1)(n-1)} \end{bmatrix} \begin{bmatrix} p_{0} \\ p_{1} \\ p_{2} \\ \vdots \\ p_{n-1} \end{bmatrix}$$

So in cases where V is as follows we have the FFT.

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Now given the values of a polynomial computed in the roots of unity how to go back to the coef representation ?

We solved $V * C_{oef} = Y$ now to solve $C_{oef} = V^{-1} * Y$. How V^{-1} looks like?

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So if $V^{-1} = \frac{1}{n}\overline{V}$ we have that the the conjugates of the roots of unity ara also roots of unity of we can apply the FFT with $\omega = \omega^{-1}$ and then divide the resulting array by 1/n.

Proof. We claim that $P = V \cdot \overline{V} = nI$:

$$p_{jk} = (\operatorname{row} j \text{ of } V) \cdot (\operatorname{col.} k \text{ of } \bar{V})$$
$$= \sum_{m=0}^{n-1} e^{ij\tau m/n} \overline{e^{ik\tau m/n}}$$
$$= \sum_{m=0}^{n-1} e^{ij\tau m/n} e^{-ik\tau m/n}$$
$$= \sum_{m=0}^{n-1} e^{i(j-k)\tau m/n}$$

Now if j = k, $p_{jk} = \sum_{m=0}^{n-1} = n$. Otherwise it forms a geometric series.

$$p_{jk} = \sum_{m=0}^{n-1} (e^{i(j-k)\tau/n})^m$$
$$= \frac{(e^{i\tau(j-k)/n})^n - 1}{e^{i\tau(j-k)/n} - 1}$$