

Mini-Course: “Games, Dynamics and Learning”

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Follow the Regularized Leader

The algorithm of Follow the Regularized Leader is defined by the round-by-round recursive rule

$$\begin{aligned} X_{i,n} &= Q_i(Y_{i,n}) \\ Y_{i,n+1} &= Y_{i,n} + \gamma_n \hat{V}_{i,n} \end{aligned} \quad (\text{FTRL})$$

- | $Q_i: Y_i \rightarrow X_i$ denotes the “choice map” of player $i \in N$.
- | $\gamma_n > 0$ is a “learning rate” parameter such that $\sum_n \gamma_n = 1$.
- | $\hat{V}_{i,n}$ is a “payoff signal” that provides an estimate for the mixed payoffs of player i at stage n .

Regularization

The second component of FTRL is the choice map

$$Q_i(y_i) = \arg \max_{x_i \in X_i} \{ y_i \cdot x_i - h_i(x_i) \}.$$

In the above, each player's *regularizer* $h_i: X_i \rightarrow \mathbb{R}$ is defined as $h_i(x_i) = \sum_{i \in A_i} \theta_i(x_i)$ for some “kernel function” $\theta_i: [0, 1] \rightarrow \mathbb{R}$ with the following properties:

- (i) θ_i is *continuous* on $[0, 1]$;
- (ii) C^2 -smooth on $(0, 1]$; and
- (iii) $\inf_{[0,1]} \theta_i > 0$.

Examples

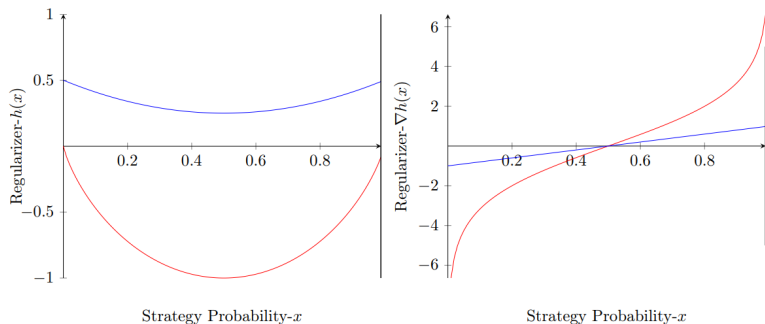
- | Negative Shannon Entropy: $h(x) = \sum_i x_i \log(x_i)$
 - | Exponential/Multiplicative Weight Updates

$$x_i(y) = \exp(y_i) / \sum_j \exp(y_j)$$

- | Euclidean Regularizer: $h(x) = \sum_i x_i^2 / 2$
 - | Euclidean Projection

$$x(y) = \arg \min_x \|y - x\|^2$$

Dichotomy of regularizers



$$\left[\begin{array}{l} \text{steep} \\ \text{non-steep} \end{array} \right. \quad \begin{array}{l} h_1(x) = x \log(x) + (1-x) \log(1-x) \\ h_2(x) = \frac{1}{2}x^2 + \frac{1}{2}(1-x)^2 \end{array} \right]$$

The feedback model

We assume a “black-box” model for players’ payoff vector of the form

$$\hat{v}_n = v(X_n) + Z_n \quad (1)$$

for some abstract error process $Z_n = (Z_{i,n})_{i \in N}$.

We will further decompose Z_n as $Z_n = U_n + b_n$, where

- | Random (zero-mean) error: $E[U_n / F_n] = 0$.
- | Systematic error: $b_n = E[Z_n / F_n]$.

with F_n denoting the history of X_n up to stage n (inclusive).

Assumptions

We may then characterize the input signal \hat{v}_n by means of the following statistics:

1. *Bias*: $E[b_n / F_n] = B_n$
2. *Variance*: $E[U_n^2 / F_n] = M_n^2$

In the above, B_n and M_n represent deterministic bounds on the bias and variance of the feedback signal \hat{v}_n .

Assumptions

For concreteness, we will also make the following blanket assumptions:

1. *Bias control*: $\lim_n B_n = 0$ and $\sum_n \gamma_n B_n < \infty$.
2. *Variance control*: $\sum_n \gamma_n^2 M_n^2 < \infty$.
3. *Generic observation errors at equilibrium*: For every mixed Nash equilibrium x of Γ and for all $n = 0, 1, \dots$, there exists a player $i \in N$ and strategies $a, b \in \text{supp}(x_i)$ such that

$$P(|\hat{v}_{ia,n} - \hat{v}_{ib,n}| \leq \beta / F_n) > 0 \quad \text{for all sufficiently small } \beta > 0.$$

Examples



Model 1 - Oracle based feedback

- | At each round n , every player $i \in N$ picks an action $\alpha_{i,n} \in A_i$ based on $X_{i,n} = X_i$.

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- | An oracle reveals to each player the pure payoff vector $v_i(\alpha_n) = (u_i(\alpha_i; \alpha_{-i,n}))_{i \in A_i}$.
- | Then the player's feedback signal is $\hat{v}_{i,n} = v_i(\alpha_n)$.

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| Assumption for bias is trivial because

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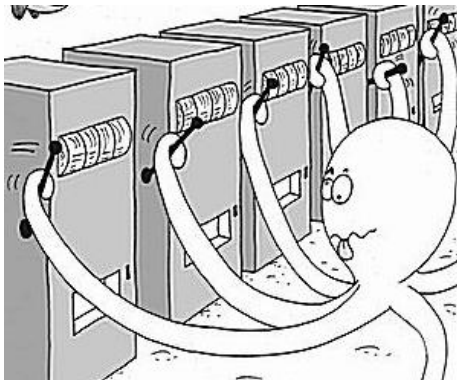
- | Assumption for bias is trivial because $E[\hat{V}_n / F_n] = E_{X_n}[v(\alpha_n)] = v(X_n)$, i.e., $b_n = 0$.
- | Assumption for noise is satisfied as long as $\sum_n \gamma_n^2 < \infty$, since $U_n \leq 2 \max_X v(X)$.

Model 1 - Oracle based feedback

Special case of our general model with

- | (A1) is trivial because $E[\hat{V}_n / F_n] = E_{X_n}[V(\alpha_n)] = V(X_n)$,
i.e., $b_n = 0$.
- | (A2) is satisfied as long as $\sum_n \gamma_n^2 < U_n$, since
 $2 \max_X V(X)$.
- | (A3) is an immediate consequence of genericity. Otherwise,
the game should have pure Nash equilibria.

Model 2 - Payoff based feedback (Bandit)



Bandit Case

- | At each round n , every player $i \in \mathcal{N}$ picks an action $\alpha_{i,n} \in A_i$ based on $X_{i,n} = X_i$.

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- | At each round n , every player $i \in N$ picks an action $\alpha_{i,n} \in A_i$ based on $X_{i,n} = X_i$.
- | Players observe their realized payoffs $u_i(\alpha_{i,n}, \alpha_{-i,n})$
- | Players need to somehow estimate their payoffs!

Importance Weighted Estimator

$$\hat{v}_{i,a,n} = \begin{cases} 0 & , \text{ if } a = a_{i,n} \\ \frac{u_i(a; a_{-i,n})}{x_{i,a,n}} & , \text{ if } a \neq a_{i,n} \end{cases}$$

- | Unbiased: $E[\hat{v}_{i,n}] = v_i(X_n)$
- | Unbounded Variance: $E[\hat{v}_{i,n}^2 / F_n] = \frac{1}{\min x_{i,a,n}}$

FTRL-exploration

- Idea: We do not limit from the beginning other options, we regularize the probabilities with an exploitation parameter that goes to zero in the infinity.

$$Y_{ia,n+1} = Y_{ia,n} + \gamma_n \hat{V}_{ia,n}$$
$$X_{i,n} = \arg \max_{X \in \Delta(A_i)} \{ Y_{i,n}, X - h_i(X) \}$$
$$\hat{X}_{i,n} = (1 - \epsilon_n) X_{i,n} + \frac{\epsilon_n}{A_j}$$

- Unbiased: $E[\hat{V}_{i,n}] = v_i(\hat{X}_n)$
- Bounded Variance: $E[U_{i,n}^2 / F_n] = \frac{1}{\min \hat{X}_{ia,n}} = O(1/\epsilon_n)$
- Bias: $b_n = v(\hat{X}_n) - v(X_n) = O(\epsilon_n)$

The Bandit Case

- | (A1) is satisfied as long as $\varepsilon_n \rightarrow 0$ and $\sum_n \gamma_n \varepsilon_n < \infty$.
- | (A2) is satisfied $\sum_n \gamma_n^2 \varepsilon_n^{-1} < \infty$.
- | (A3) is an immediate consequence of genericity. Otherwise, the game should have pure Nash equilibria.

Asymptotic Stability

A point $x \in X$ is said to be

1. *Stochastically stable* under (FTRL): If for all $\delta > 0$ and all neighborhoods U of x there exists open set of initial conditions $W_0 \subset Y$ such that

$$P(X_n \in U \text{ for all } n = 0, 1, \dots) \geq 1 - \delta$$

whenever $Y_0 \subset W_0$.

2. *Stochastically attracting* under (FTRL): If for all $\delta > 0$, there exists open set of initial conditions $W_0 \subset Y$ such that

$$P(\lim_n X_n = x) \geq 1 - \delta$$

whenever $Y_0 \subset W_0$.

3. *Stochastically asymptotically stable* under (FTRL): if it is stochastically stable and attracting.

Main Results

Main Theorem. Suppose that Assumptions 1–3 hold.

Then:

x is a strict Nash equilibrium x is stochastically
asymptotically stable under (FTRL)

Main Results

Theorem

Let $x^ \in X$ be a strict Nash equilibrium of Γ . If (FTRL) is run with inexact payoff feedback satisfying Assumptions 1 and 2, then x^* is stochastically asymptotically stable.*

Theorem

Let x^ be a mixed Nash equilibrium of Γ . If (FTRL) is run with inexact payoff feedback satisfying assumption 3, then x^* is not stochastically asymptotically stable.*

Proof techniques - Instability

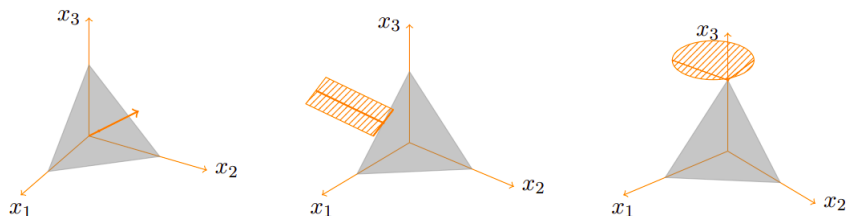


Figure: Polar cone

1. $x = Q(y) \quad y \in \partial h(x)$
2. $\partial h(x) = h(x) + PC(x)$ for all $x \in X$,
 where $PC(x) = \{y \in Y : y_a \leq y_b \text{ for all } a, b \in A\}$.

Proof techniques - Instability

Lemma (Informal)

Let $X_{i,n}$ be the sequence of play in (FTRL) i.e., $X_{i,n} = Q(Y_{i,n})$
 X_i of player $i \in N$; and for some round $n \geq 0$ let $a, b \in \text{supp}(X_{i,n})$
be two pure strategies of player $i \in N$. Then it holds:

$$(\theta_i(X_{ia,n+1}) - \theta_i(X_{ia,n})) - (\theta_i(X_{ib,n+1}) - \theta_i(X_{ib,n})) = \gamma_n(\hat{v}_{ia,n} - \hat{v}_{ib,n})$$

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- | Assume ad absurdum that a mixed Nash equilibrium x is stochastically asymptotically stable. Since x is mixed, there exist $a, b \in \text{supp}(x)$.

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$$(\theta_i(X_{ia,n+1}) - \theta_i(X_{ia,n})) - (\theta_i(X_{ib,n+1}) - \theta_i(X_{ib,n})) = \gamma_n(\hat{v}_{ia,n} - \hat{v}_{ib,n})$$

- | Assume ad absurdum that a mixed Nash equilibrium x is stochastically asymptotically stable. Since x is mixed, there exist $a, b \in \text{supp}(x)$.
- | The stochastic stability implies that for all $\varepsilon, \delta > 0$ if X_0 belongs to an initial neighborhood U , then $\|X_n - x\| < \varepsilon$ for all $n \geq 0$, with probability at least $1 - \delta$.

Proof techniques - Instability

- | By the triangle inequality for two consecutive instances of the sequence of play $X_{i,n}, X_{i,n+1}$ for any player $i \in N$ it holds:

$$|X_{ia,n+1} - X_{ia,n}| + |X_{ib,n+1} - X_{ib,n}| < O(\varepsilon) \text{ with probability } 1 - \delta$$

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- | Consider ε sufficiently small, such that the probabilities of the strategies that belong to the support of the equilibrium are bounded away from 0, for all the points of the neighborhood. Since θ_i is continuously differentiable in $(0, 1]$, the differences described in the lemma above are bounded from $O(\varepsilon)$.

Proof techniques - Instability

- | If the sequence of play X_n is contained to an ε -neighborhood of x^* , then the difference of the feedback, for any player $i \in N$, to two strategies of the equilibrium is $O(\varepsilon/\gamma_n)$ with probability at least $1 - \delta$:

$$P(|\hat{v}_{ia,n} - \hat{v}_{ib,n}| = O(\varepsilon/\gamma_n) \mid F_n) \geq 1 - \delta$$

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- | From assumption 3 for a fixed round n and some player $i \in N$, there exist $\beta, \pi > 0$ such that:
 $P(|\hat{v}_{ia,n} - \hat{v}_{ib,n}| \geq \beta \mid F_n) = \pi > 0.$

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- | From assumption 3 for a fixed round n and some player $i \in N$, there exist $\beta, \pi > 0$ such that:
 $P(|\hat{v}_{ia,n} - \hat{v}_{ib,n}| \leq \beta \mid F_n) = \pi > 0$.
- | Thus by choosing $\varepsilon = O(\beta\gamma_n)$ and $\delta = \pi/2$, we obtain a contradiction and our proof is complete.

Nash equilibria - reminder

A point x is a *Nash equilibrium* of Γ if

$$u_i(x) \geq u_i(x_i; x_{-i}) \quad \text{for all } x_i \in X_i \text{ and all } i \in N. \quad (\text{NE})$$

We call support of x the set: $\text{supp}(x) = \{\alpha_i \in A_i : x_{\alpha_i} > 0\}$.
Equivalently, Nash equilibria can be characterized by means of the variational inequality

$$v_{\alpha_i}(x) \leq v_{\alpha_i}(x) \quad \text{for all } \alpha_i \in \text{supp}(x) \text{ and all } \alpha_i \in A_i, i \in N.$$

Proof techniques - Stability

- | Let $x = (\alpha_1, \dots, \alpha_N) \in A$ be a strict Nash equilibrium. Then for every $\varepsilon \in (0, 1)$, there exist constants M_i and the corresponding score-dominant open sets for each player $i \in N$ such that: $\prod_{i \in N} Q_i(W_i(M_i, \varepsilon)) \subseteq U$, where $U = \{x \in X : x_i > 1 - \varepsilon \text{ for all } i \in N\}$ and

$$W_i(M_i) = \{Y_i : Y_{i,i} - Y_{i,j} > M_i \text{ for all } \alpha_j = \alpha_j, \alpha_j \in A_j\}$$

for each player $i \in N$

Proof techniques - Stability

- | Fix a confidence level $\delta > 0$, focus on one player $i \in N$ and drop the index i for simplicity; consider a neighborhood U of x^* that can be described as the one above and for which $u_i(X) - u_i(x^*) \geq -c$ for some $c > 0$, for all $\alpha = \alpha_i, \alpha \in A_i$ and all $X \in U$.
- | We will prove by induction that there exists an open set of initial conditions W_0 , such that whenever $Y_0 \in W_0$ then $Y_n \in W$ for all $n = 0, 1, \dots$.

- Notice that whenever $X \in U$, the payoffs belong to the set $W = Q^{-1}(U)$. Furthermore, the payoff differences $Y_i - Y_j$ between every pure strategy $\alpha \in A_i$, $\alpha \neq \alpha^*$ and the strategy of the equilibrium α^* can be expressed as

$$Y_{i,n+1} - Y_{j,n+1} = Y_{i,0} - Y_{j,0} + \sum_{k=0}^n \gamma_k (u_i(X_k) - u_j(X_k)) \\ + \sum_{k=0}^n \gamma_k \text{Noise}_k + \sum_{k=0}^n \gamma_k \text{Bias}_k$$

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$$Y_{i,n+1} - Y_{i,0} = Y_{i,0} - Y_{i,0} + \sum_{k=0}^n \gamma_k (u_i(X_k) - u_i(X_k^*)) \\ + \sum_{k=0}^n \gamma_k \text{Noise}_k + \sum_{k=0}^n \gamma_k \text{Bias}_k$$

- | Using martingale limit theory we control the terms $\sum_{k=0}^n \gamma_k \text{Noise}_k$, $\sum_{k=0}^n \gamma_k \text{Bias}_k$ as to be less than $\varepsilon_1 = \sqrt{2 \sum_{k=0}^n \gamma_k^2 M_k^2 / \delta}$, $\varepsilon_2 = 2 \sum_{k=0}^n \gamma_k B_k / \delta$ equivalently with probability at least $1 - \delta$.

- | Let $R_n = \sum_{k=0}^n \gamma_k (U_{a,k} - U_{b,k})$, which is a martingale.
- | Consider the event $D_{n, \varepsilon_1} = \{\sup_{0 \leq k \leq n} |R_k| \leq \varepsilon_1\}$, then

$$P(D_{n, \varepsilon_1}) \geq \frac{E[R_n^2]}{\varepsilon_1^2} = \frac{2 \sum_{k=0}^n \gamma_k^2 M_k^2}{\varepsilon_1^2}$$

- | Notice that

$$\begin{aligned} E[R_n^2] &= \sum_{k=0}^n \gamma_k^2 E[(U_{a,k} - U_{b,k})^2] = 2 \sum_{k=0}^n \gamma_k^2 E[U_k^2] \\ &= 2 \sum_{k=0}^n \gamma_k^2 E[E[U_k^2 | F_k]] = 2 \sum_{k=0}^n \gamma_k^2 M_k^2 \end{aligned}$$

and $E[U_{a,k} U_{b,l}] = E[E[U_{a,k} U_{b,l} | F_{k+l}]] = 0$ for all $k \neq l$ and a, b be either of the pure strategy α and the strategy of the equilibrium α^* , due to the noise being zero-mean.

- | Let $\Gamma_1 = 2 \sum_{k=0} \gamma_k^2 M_k^2$ and choose $\varepsilon_1 = \sqrt{2\Gamma_1/\delta}$.
- | The event $D_1 = \bigcap_{n=0} D_{1,n}$ will happen with probability at most $\delta/2$.

| Notice that

$$\left| \sum_{k=0}^n \gamma_k (b_{k+1} - b_k) \right| = \sum_{k=0}^n \gamma_k |b_{k+1} - b_k| \leq 2 \sum_{k=0}^n \gamma_k b_k$$

| Let $S_n = 2 \sum_{k=0}^n \gamma_k b_k$, which is a submartingale.

| If $E_{n,2} = \{\sup_{0 \leq k \leq n} S_k \leq \varepsilon_2\}$ then it holds

$$P(E_{n,1}) \leq \frac{E[S_n]}{\varepsilon_2} = \frac{2 \sum_{k=0}^n \gamma_k E[E[b_k | F_k]]}{\varepsilon_2} = \frac{2 \sum_{k=0}^n \gamma_k B_k}{\varepsilon_2}$$

- | Let $\Gamma_2 = 2 \sum_{k=0} \gamma_k B_k$ and choose $\varepsilon_2 = 2\Gamma_2/\delta$.
- | Then the event $E_2 = \bigcap_{n=0} E_{n, 2}$ will occur with probability at most $\delta/2$.

- I Choose $M_0 > M + \varepsilon_1 + \varepsilon_2$ and let $W_0 = \{Y : Y < -M_0 \text{ for all } \alpha = \alpha\}$. If $Y_0 \in W_0$ then with probability at least $1 - \delta$ we prove that $Y_n > M$ for all $n = 1, 2, \dots$ and thus the equilibrium is stochastically stable.

- | Choose $M_0 > M + \varepsilon_1 + \varepsilon_2$ and let $W_0 = \{Y : Y < -M_0 \text{ for all } \alpha = \alpha\}$. If $Y_0 \in W_0$ then with probability at least $1 - \delta$ we prove that $Y_n > M$ for all $n = 1, 2, \dots$ and thus the equilibrium is stochastically stable.
- | Since with probability at least $1 - \delta$ the sequence remains in the neighborhood U we have

$$Y_{n+1} - Y_{n+1} = -C \sum_{k=0}^n \gamma_k + \varepsilon_1 + \varepsilon_2 \quad (2)$$

which implies that the score differences go to $-\infty$, thus all the strategies except for the strategy of the equilibrium become dominated. As a result the point is stochastically asymptotically stable.

Permitted parameters

The above conditions for the method's learning rate and exploration parameters can be achieved by using schedules of the form

$$\begin{aligned} | \quad \gamma_n &= 1/n^p \\ | \quad \varepsilon_n &= 1/n^q \end{aligned}$$

with $p + q > 1$ and $2p - q > 1$. A popular choice is $p = 2/3 + \delta$ and $q = 1/3 + \delta$ for some arbitrarily small $\delta > 0$ – or $\delta = 0$ and including an extra logarithmic factor.