

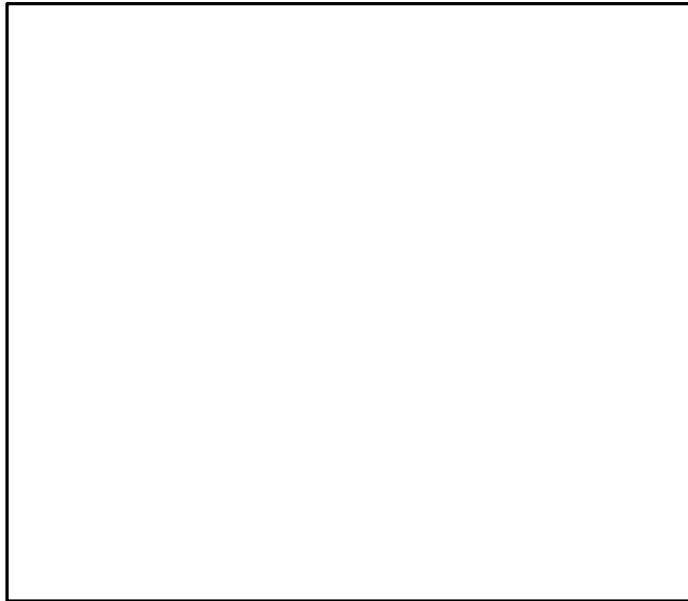


Sampling and approximation algorithms
for Gibbs point processes

Tobias Friedrich, **Andreas Göbel**, Max Katzmann, Martin Krejca, Marcus Pappik

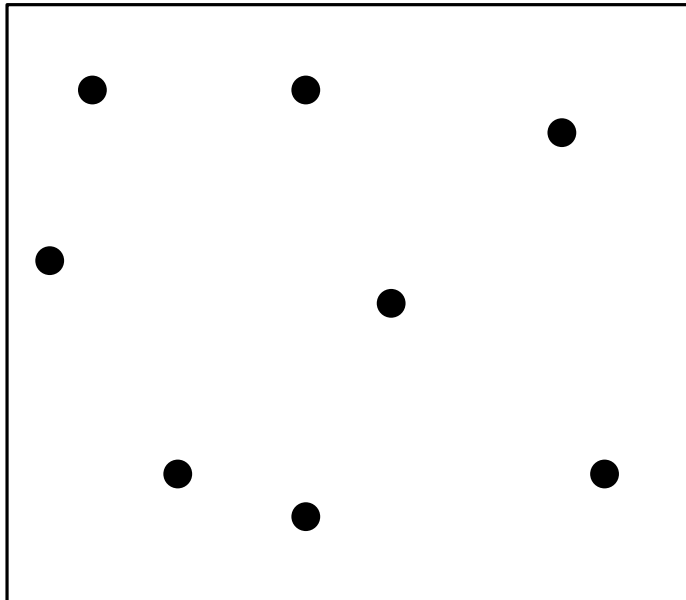
distribution of particles with *hard constraints* in (Euclidean) space

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$V \subset \mathbb{R}^d$ bounded, measurable

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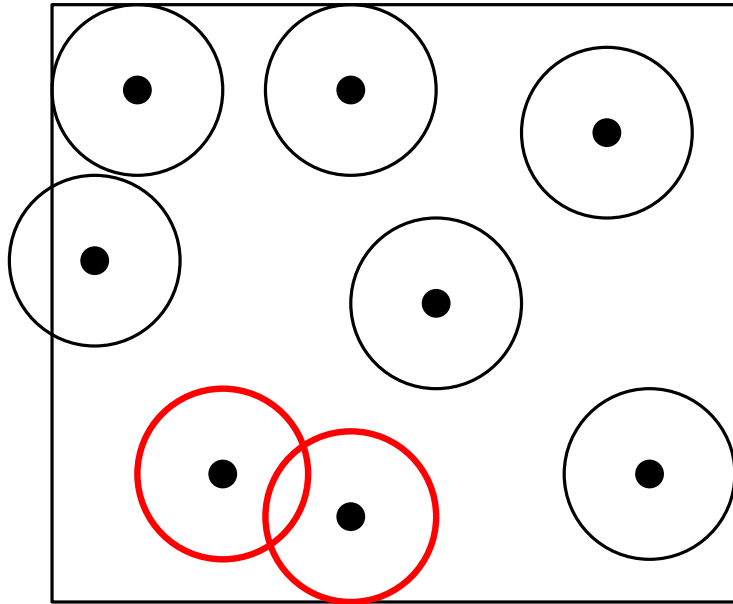


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particle centers:

distributed according to Poisson point process of intensity $\lambda \in \mathbb{R}_{\geq 0}$ (*fugacity*) on V

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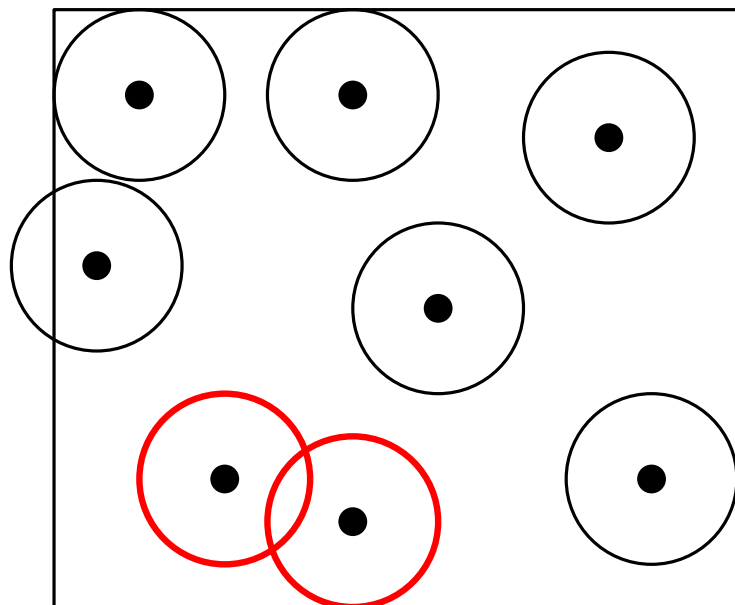
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Gibbs measure:

$$\frac{dP_{\mathbb{V}}^{(\lambda, r)}}{dQ_{\lambda}}(x_1, \dots, x_k) = \frac{D_r(x_1, \dots, x_k) e^{\lambda v(\mathbb{V})}}{\Xi_{\mathbb{V}}(\lambda, \phi)}$$

Partition function:

$$\Xi_{\mathbb{V}}(\lambda, r) = 1 + \sum_{k \in \mathbb{N}_{\geq 1}} \frac{\lambda^k}{k!} \int_{\mathbb{V}^k} D_r(x_1, \dots, x_k) v^k(d\mathbf{x})$$

Results:

- Metropolis et al. 1953
- non-rigorous results (Wilfred et al. 1998/2000, Moka et al. 2018)
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Rigorous run-time guarantees: run-time polynomial in $\nu(\mathbb{V})$

Can we  approximately sample
from the Gibbs distribution

Can we  compute the partition function
approximate

For which parameter range

Phase Transitions:

- Meyer in the 40's: no phase transition for $\lambda < \frac{1}{e\nu(B_{2r})}$
 - Cluster expansion convergence

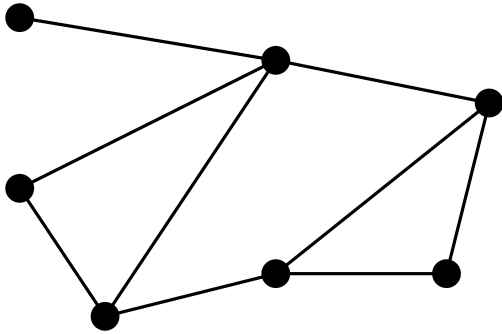
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Recent breakthroughs by translating CS methods used on discrete spin systems.

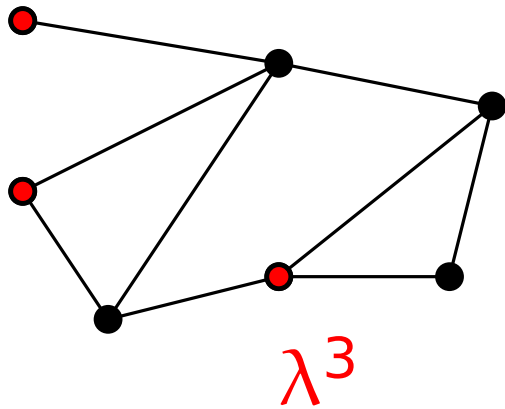
- Helmuth et al. 2020: no phase transition for $\lambda < \frac{2}{\nu(B_{2r})}$
 - Decay of correlations
- Michelen et al. 2020: no phase transition for $\lambda < \frac{e}{\nu(B_{2r})}$
 - Analyticity of the pressure
- Michelen et al. 2021: no phase transition for $\lambda < \frac{e}{(1-8^{-d-1})\nu(B_{2r})}$
 - Potential weighted constant

undirected graph $G = (V, E)$ and parameter $\lambda \in \mathbb{R}_{\geq 0}$

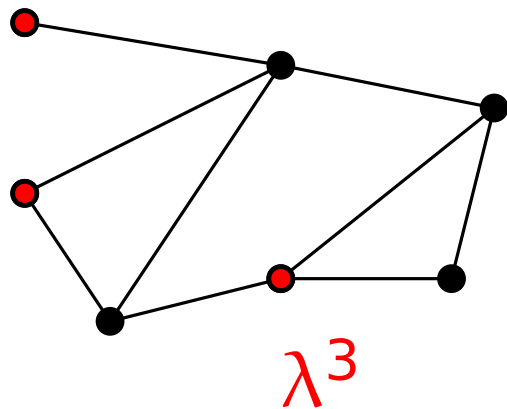


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independent set $I \in \mathcal{I}(G)$ has weight $\lambda^{|I|}$



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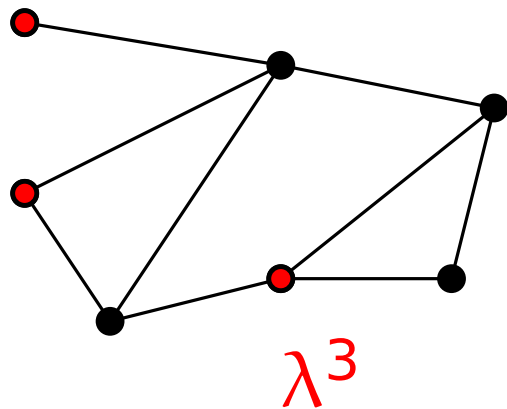


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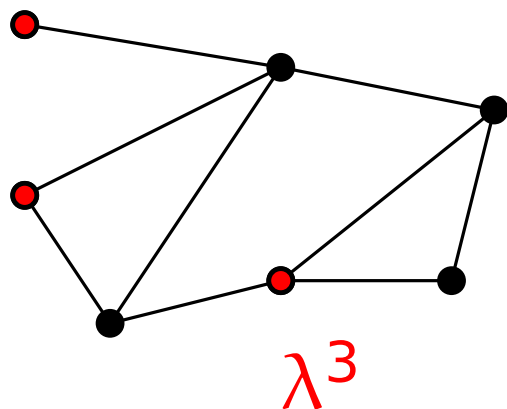
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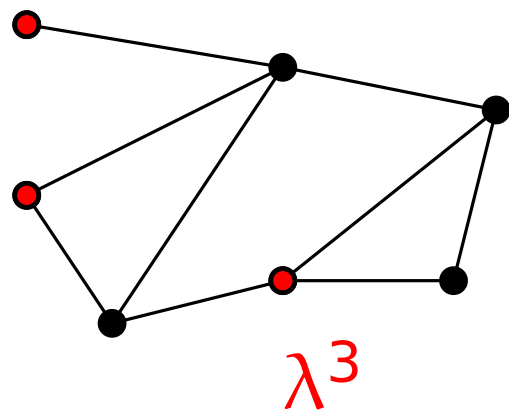
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(Weitz 2006, Barvinok 2016, Anari et al. 2020 → *Anari et al. 2021)

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NP-hard to approximate if $\lambda > \lambda_c(\Delta)$
(Sly 2010, Galanis et al. 2011)

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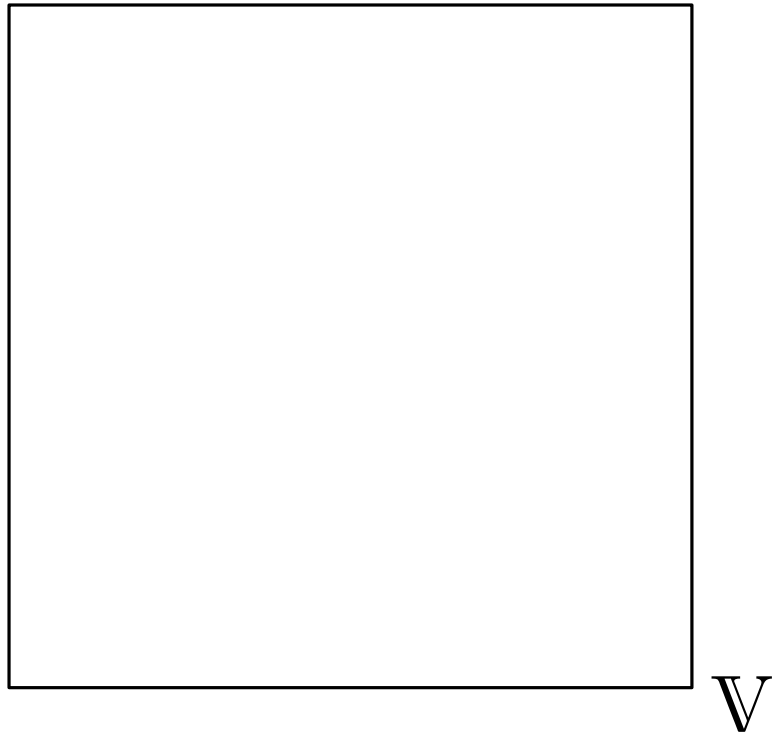


approx. $Z(G_\rho, \lambda_\rho)$

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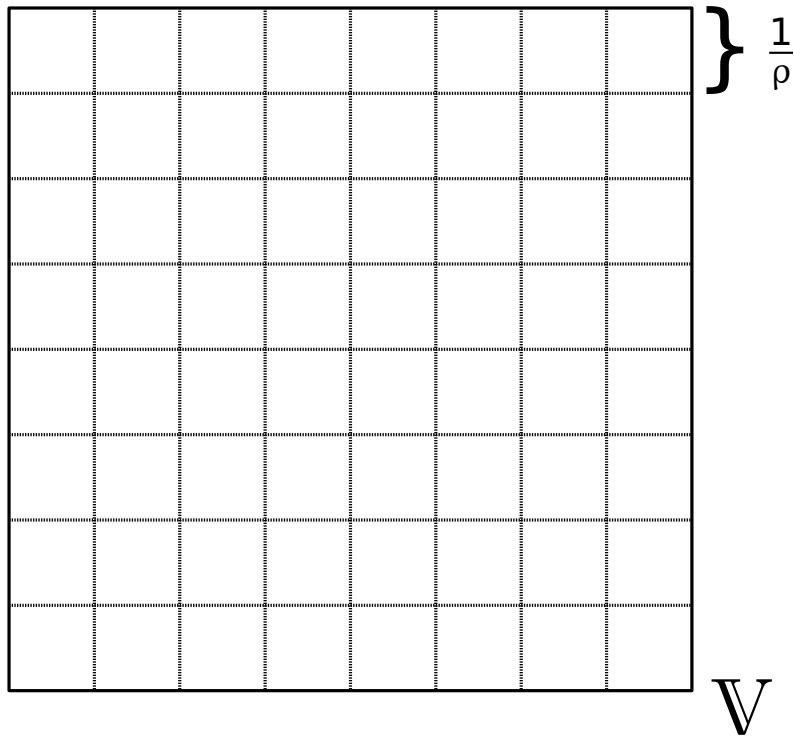
Result is also approximation for $\Xi_{\mathbb{V}}(\lambda, r)$!

1st result: Discretization



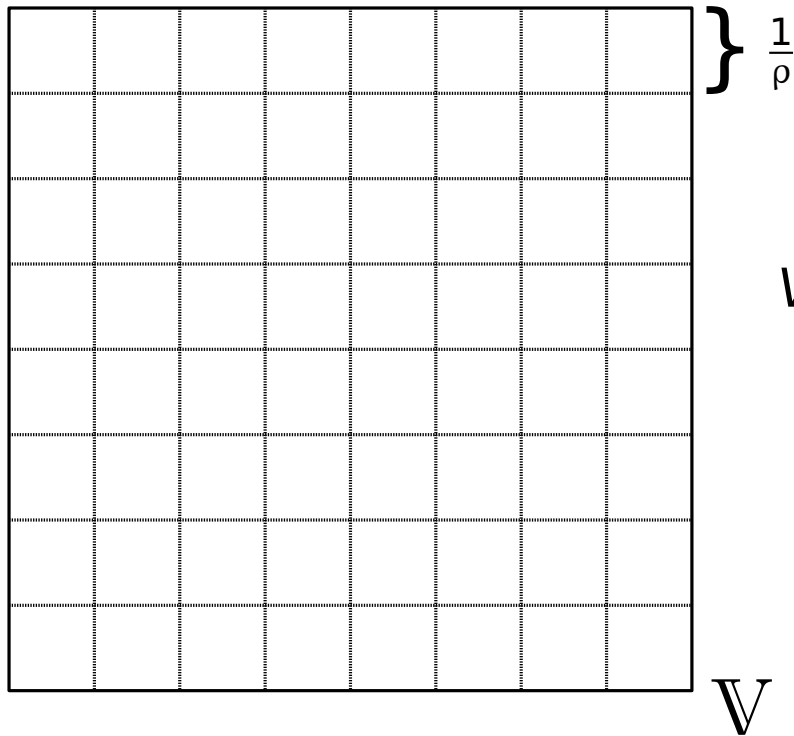
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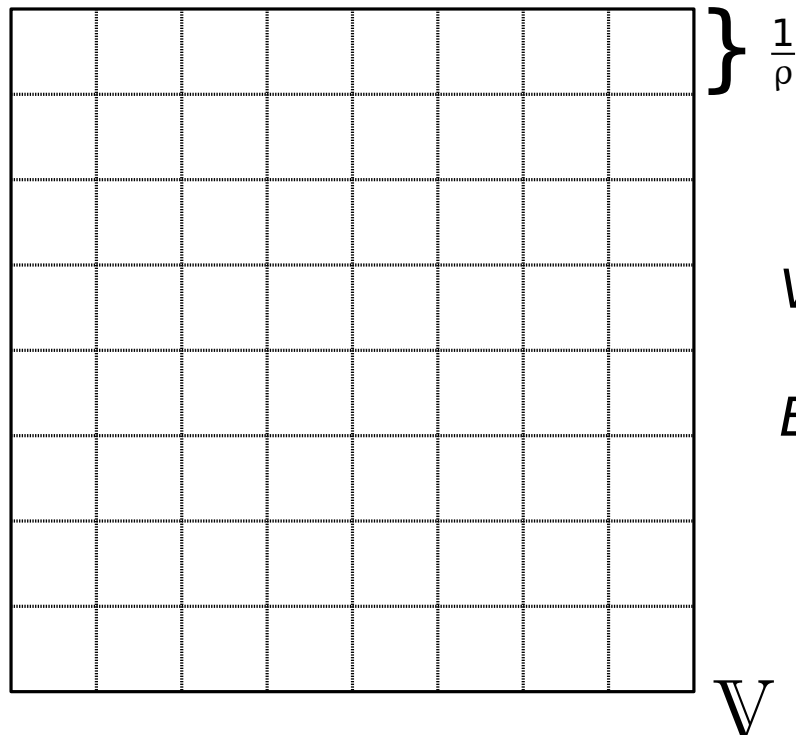
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V_ρ : contains vertex v_x for each grid point x

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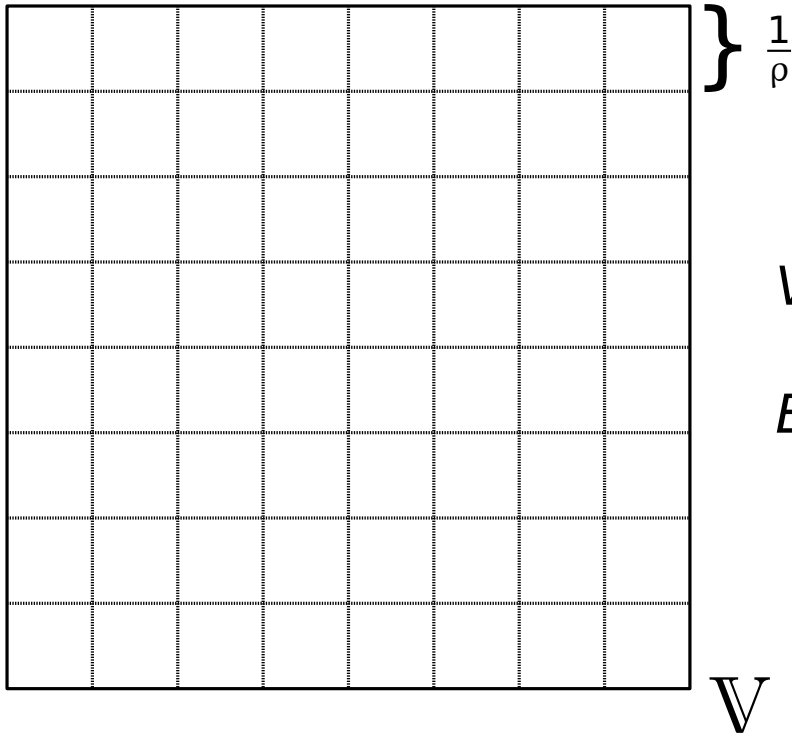


resolution $\rho \in \mathbb{R}_{\geq 1}$

V_ρ : contains vertex v_x for each grid point x

E_ρ : edge between v_x, v_y iff $x \neq y$ and $d(x, y) < 2r$

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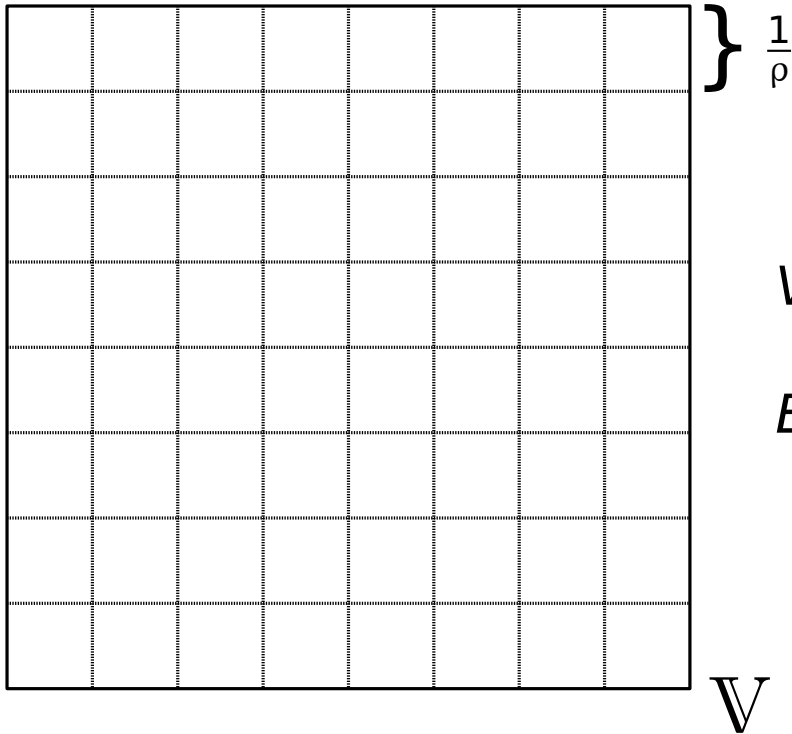
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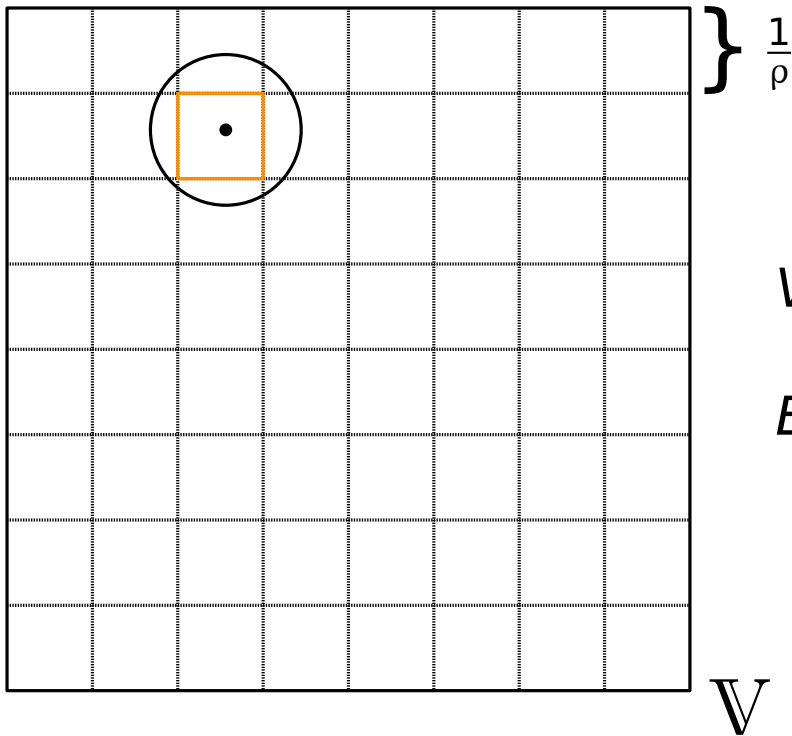
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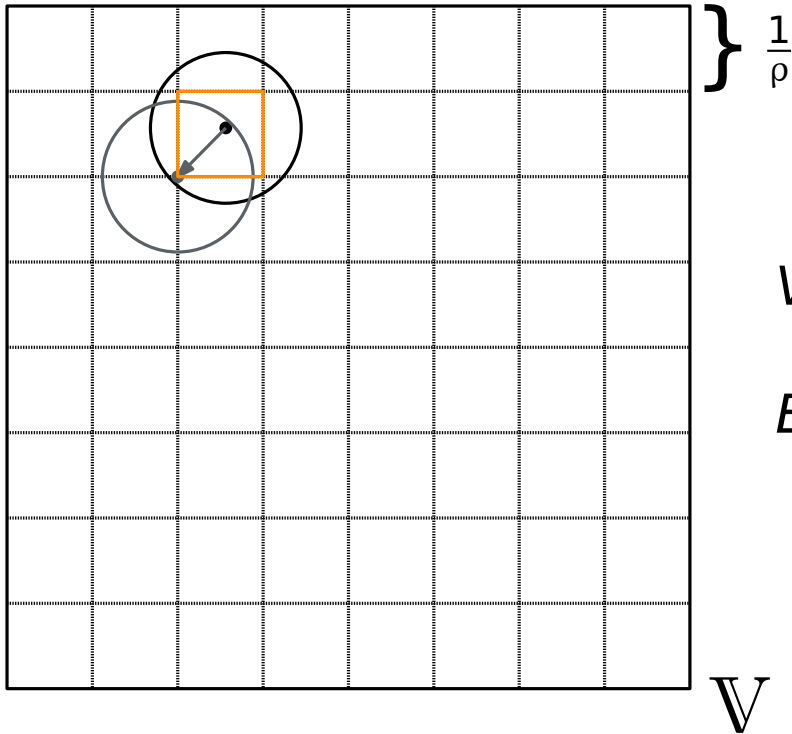
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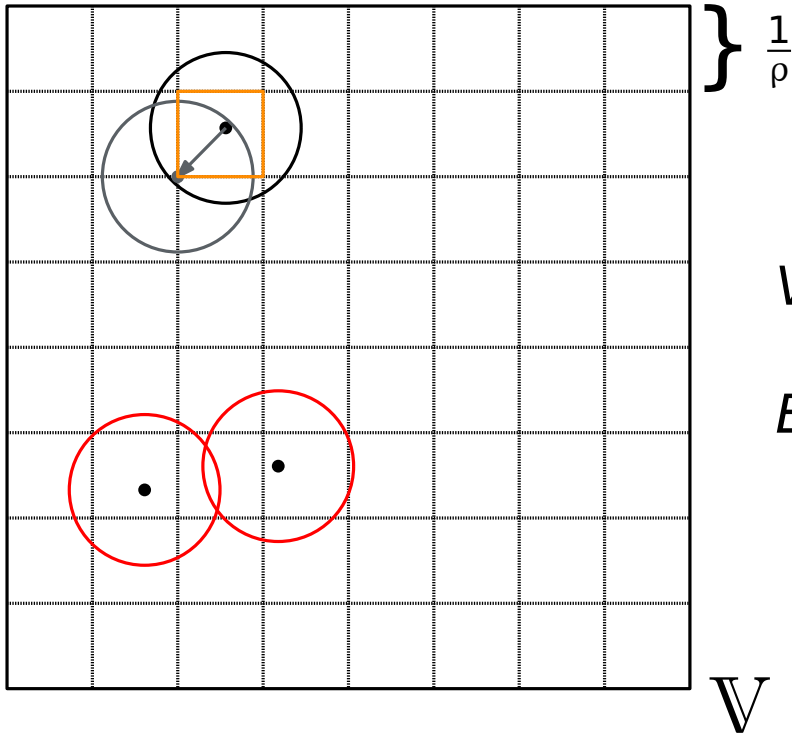
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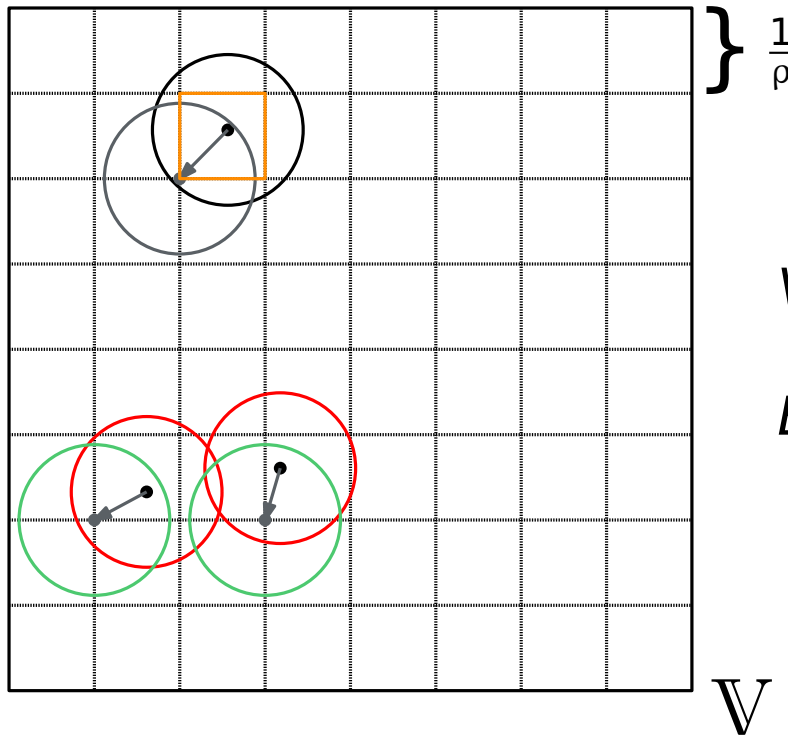
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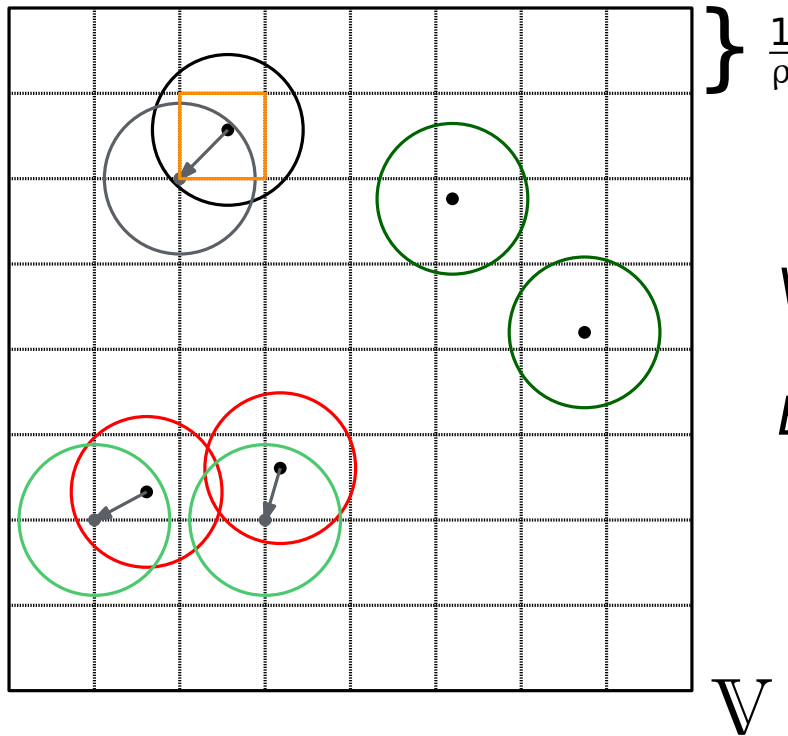
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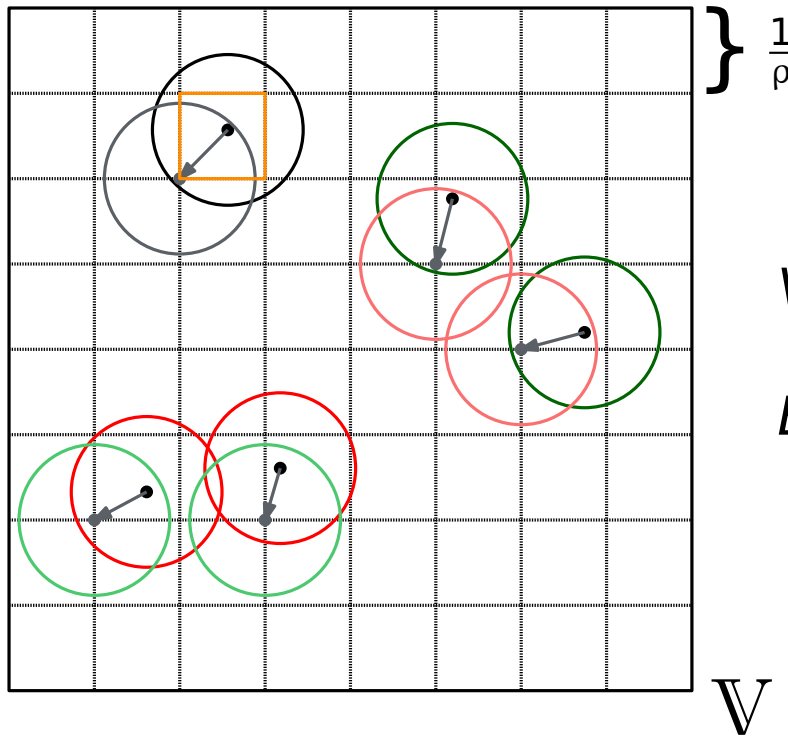
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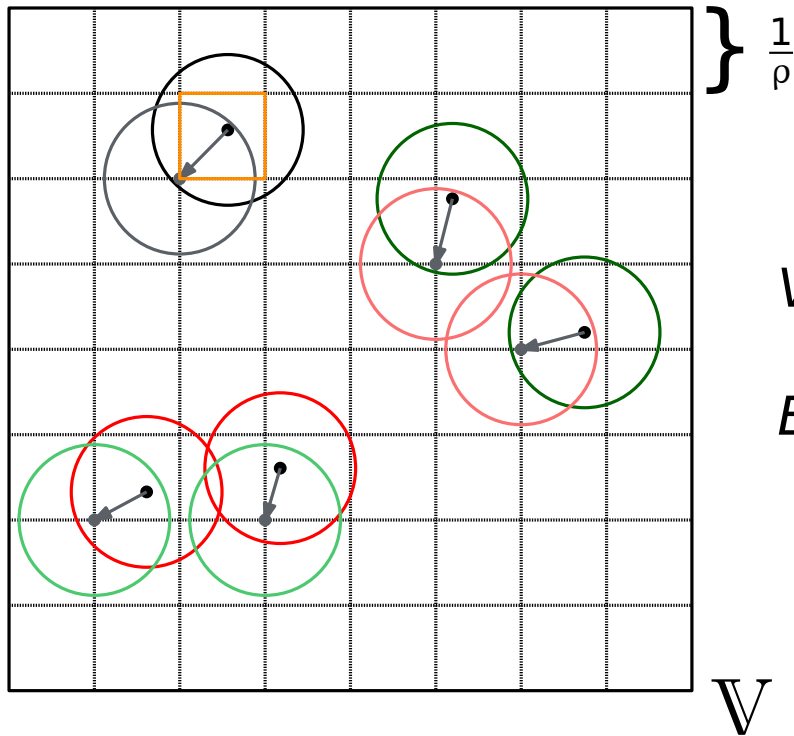
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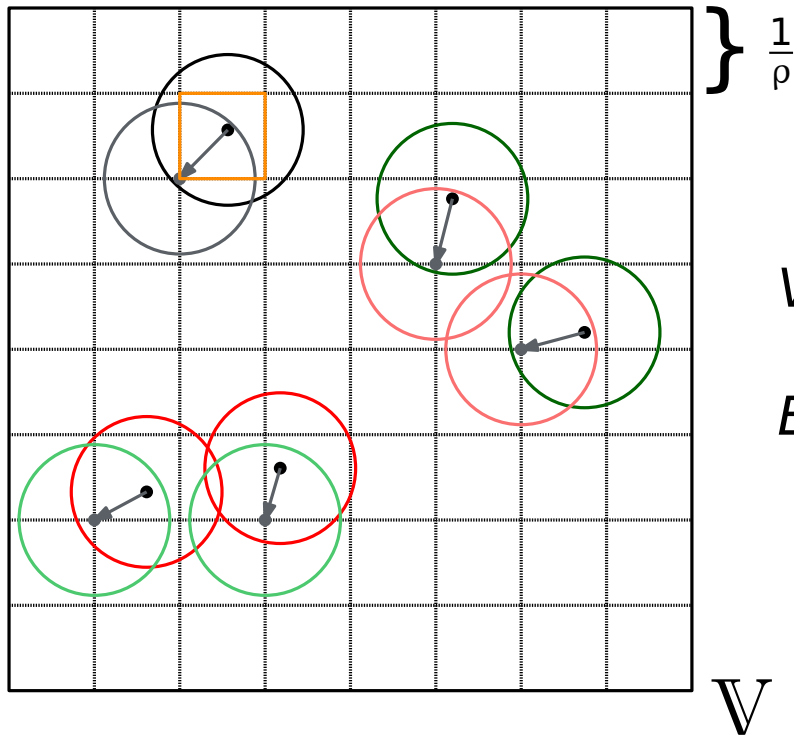
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Intuition: $\lim_{\rho \rightarrow \infty} Z(G_\rho, \lambda_\rho) = \Xi_V(\lambda, r)$

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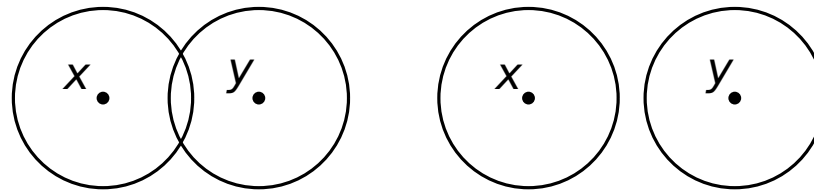
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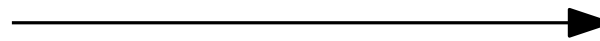
1st Result: Bounding the volume of erratic configs

Observation: discretization 'moves' particle centers by at most $\Theta\left(\frac{1}{\rho}\right)$

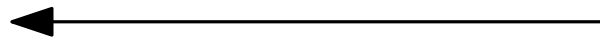
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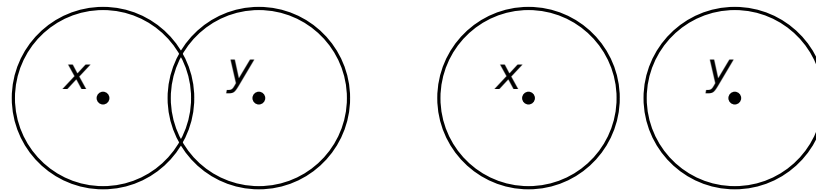
invalid to valid



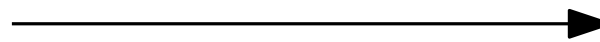
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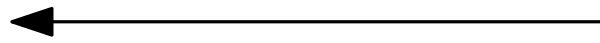


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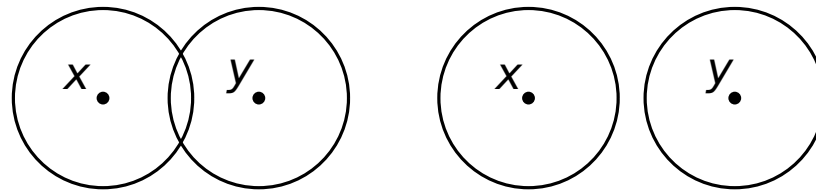


$$2r - \Theta\left(\frac{1}{\rho}\right) \leq d(x, y) < 2r$$

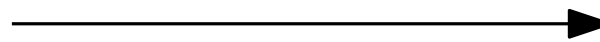
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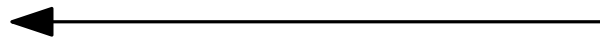


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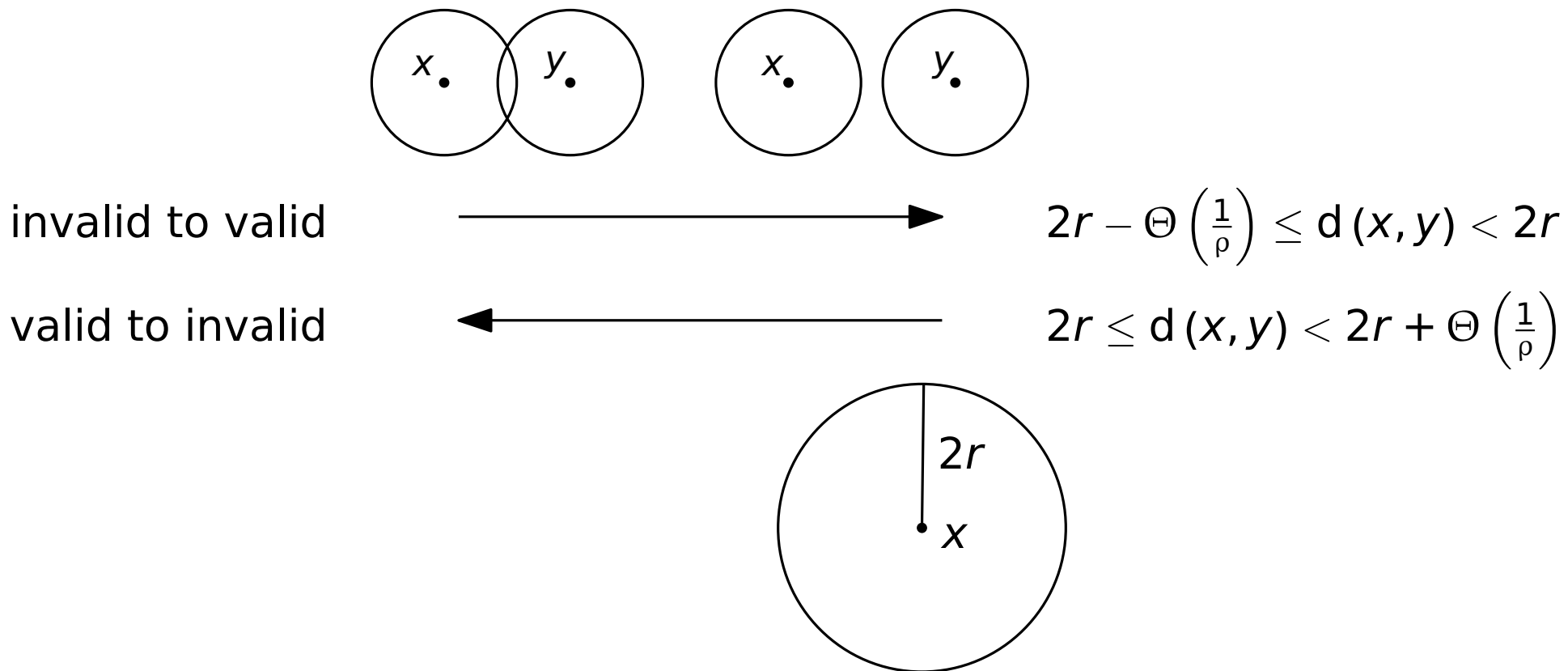
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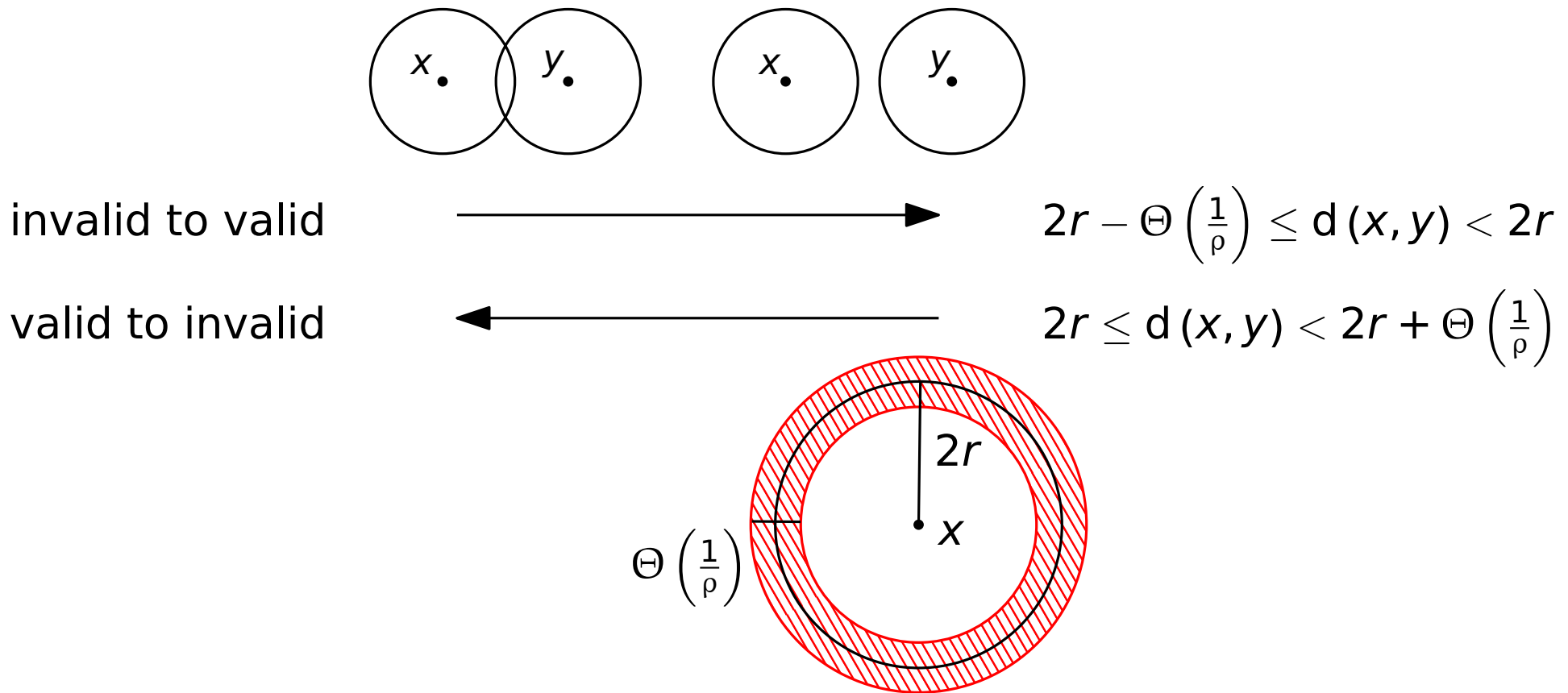


$$2r \leq d(x, y) < 2r + \Theta\left(\frac{1}{\rho}\right)$$

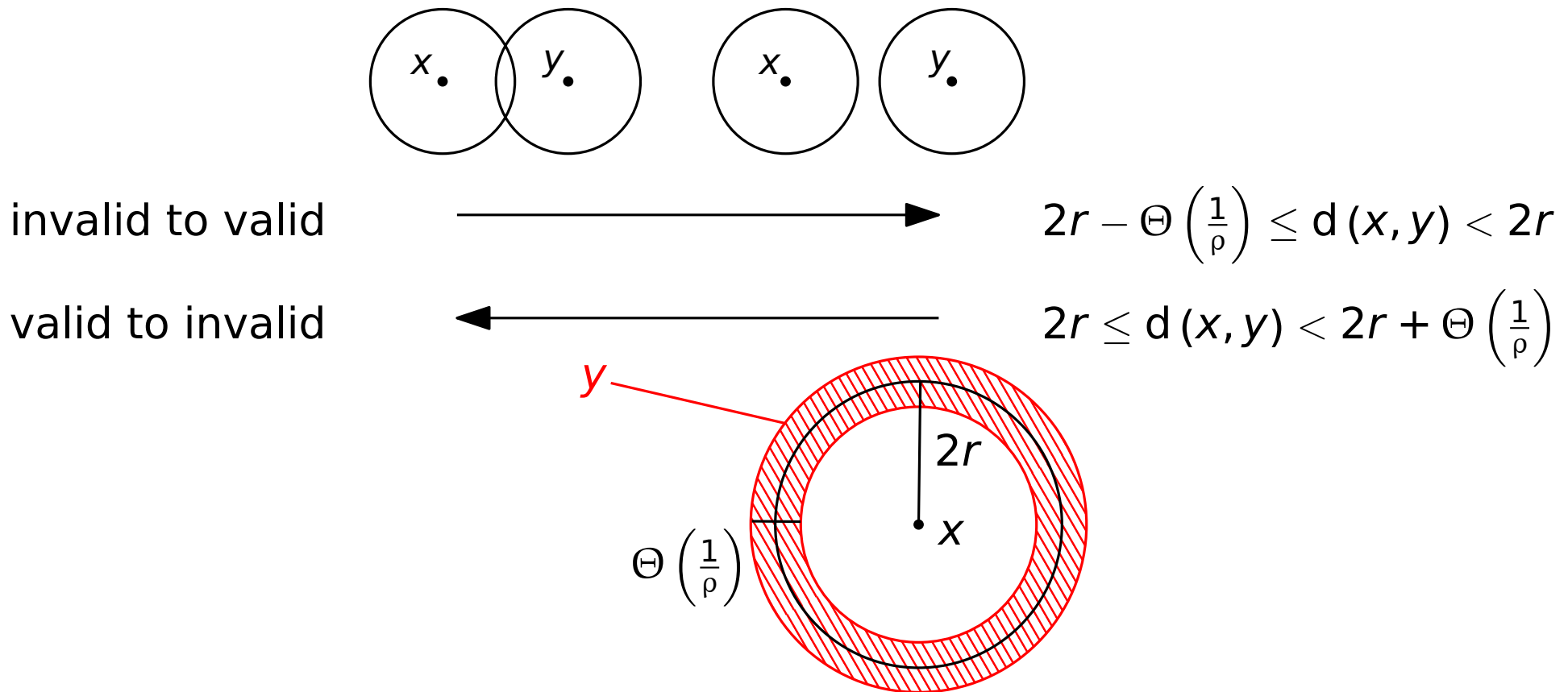
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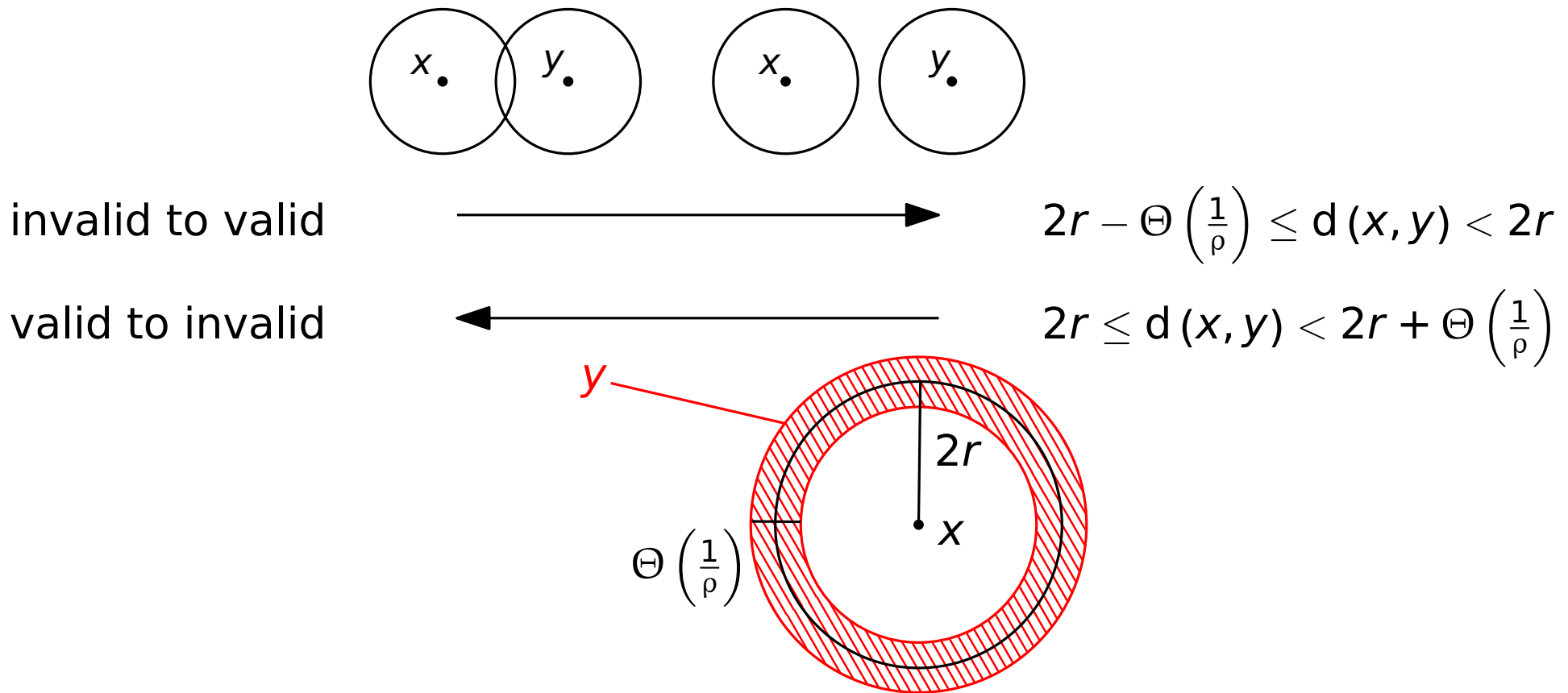
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Theorem: $|\Xi_{\mathbb{V}}(\lambda, r) - Z(G_{\rho}, \lambda_{\rho})| \leq \frac{\exp(\text{vol}(\mathbb{V}) \ln(\text{vol}(\mathbb{V})))}{\rho} \cdot \Xi_{\mathbb{V}}(\lambda, r)$

1st Result: Determining the maximum degree

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we want:

$$\frac{\lambda}{\rho^d} = \lambda_{\rho} < \lambda_c(\Delta_{\rho}) = \frac{(\Delta_{\rho}-1)^{\Delta_{\rho}-1}}{(\Delta_{\rho}-2)^{\Delta_{\rho}}} \left(\approx \frac{e}{\Delta_{\rho}} \right)$$

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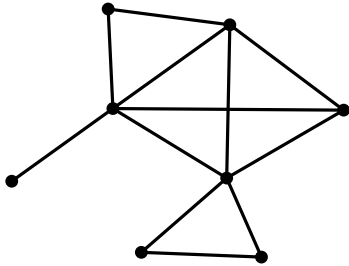
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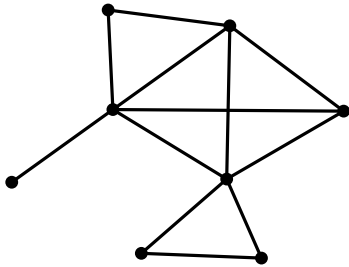
we have $|V_{\rho}| \approx \rho^d \text{vol}(\mathbb{V})$ we would need $\rho \in \Theta(\exp(\text{vol}(\mathbb{V}) \ln(\text{vol}(\mathbb{V}))))$

Existing algorithms would not run in time poly $(\text{vol}(\mathbb{V}))!$

Glauber Dynamics

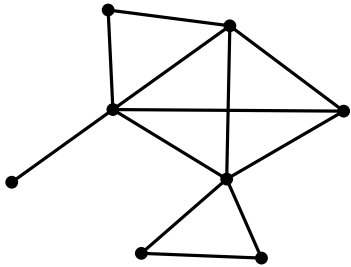


Glauber Dynamics



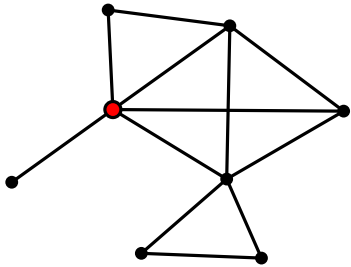
start with some (deterministic) independent set

Glauber Dynamics



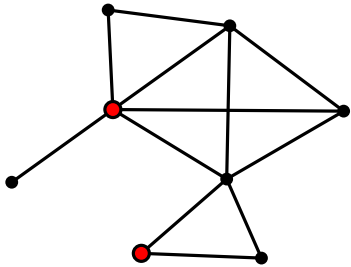
start with some (deterministic) independent set
repeat: choose vertex uniformly at random
update vertex with appropriate probability

Glauber Dynamics



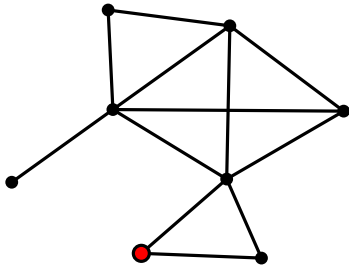
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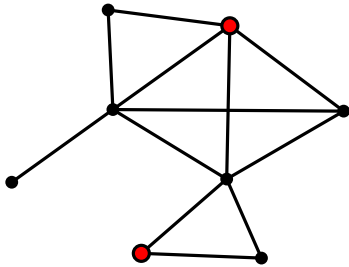
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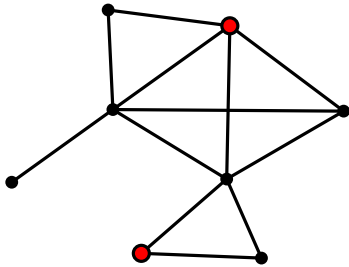
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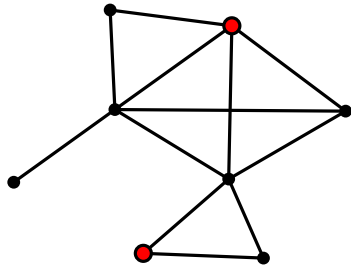
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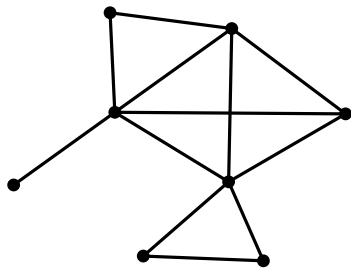
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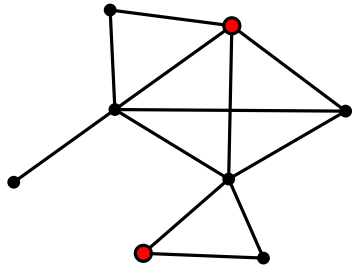


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Clique Dynamics



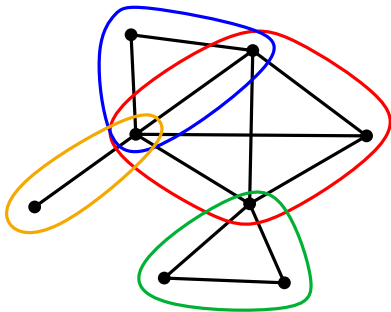
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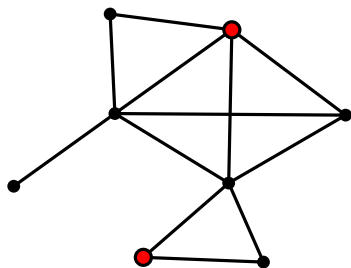
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Clique Dynamics

idea: use clique cover to update multiple vertices



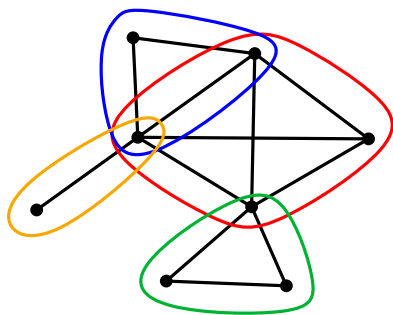
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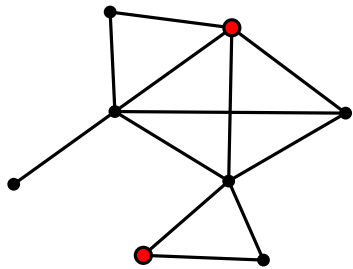
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repeat: choose *clique* uniformly at random
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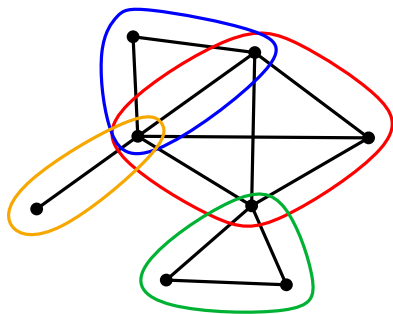
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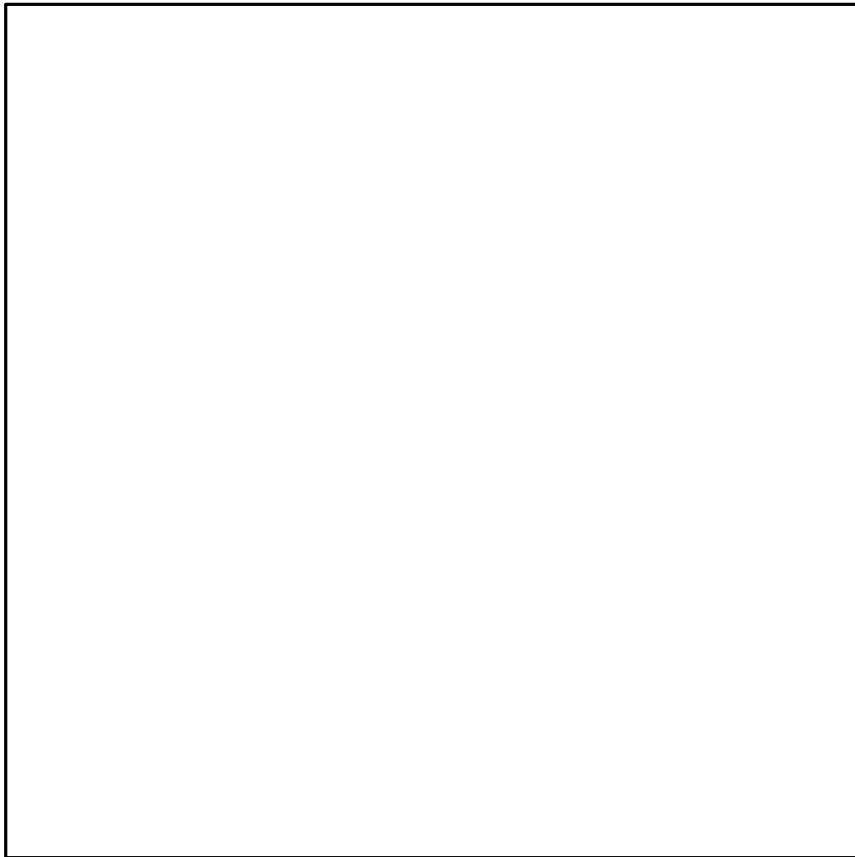


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Runtime only depends on the size of the clique cover!

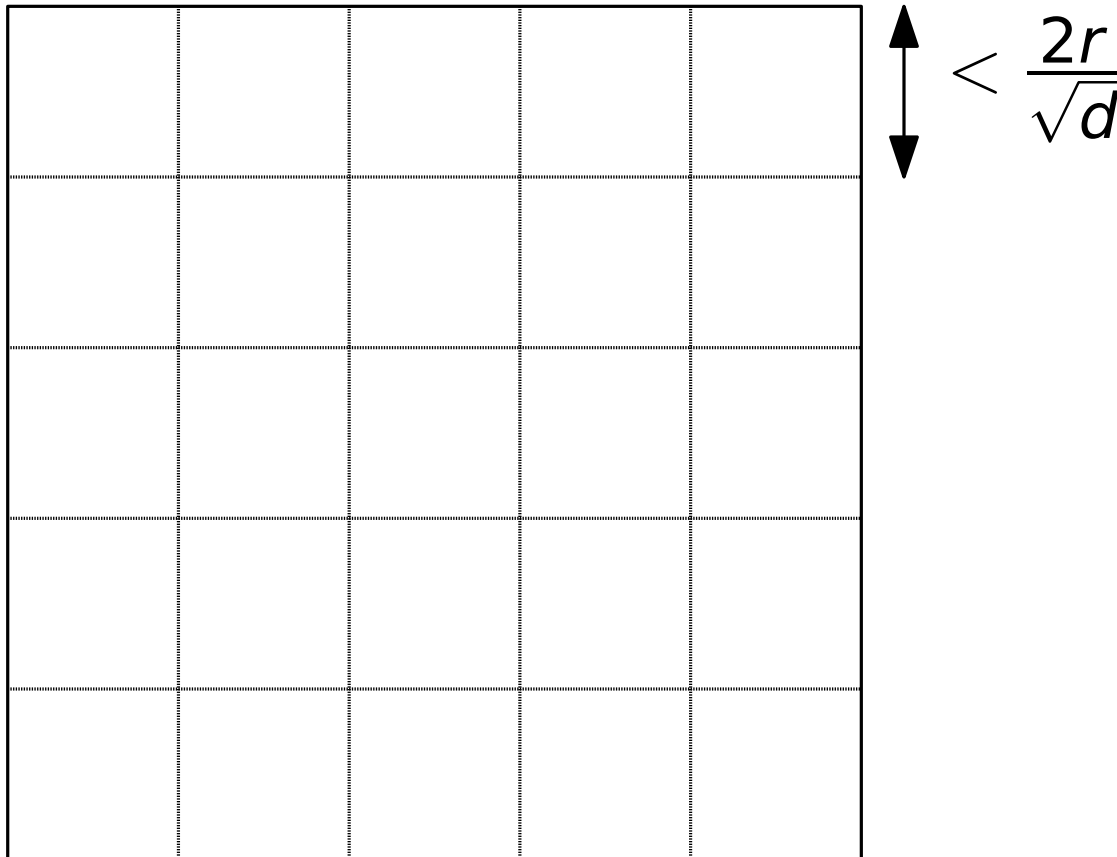
1st Result: Cliques in our Discretization

V



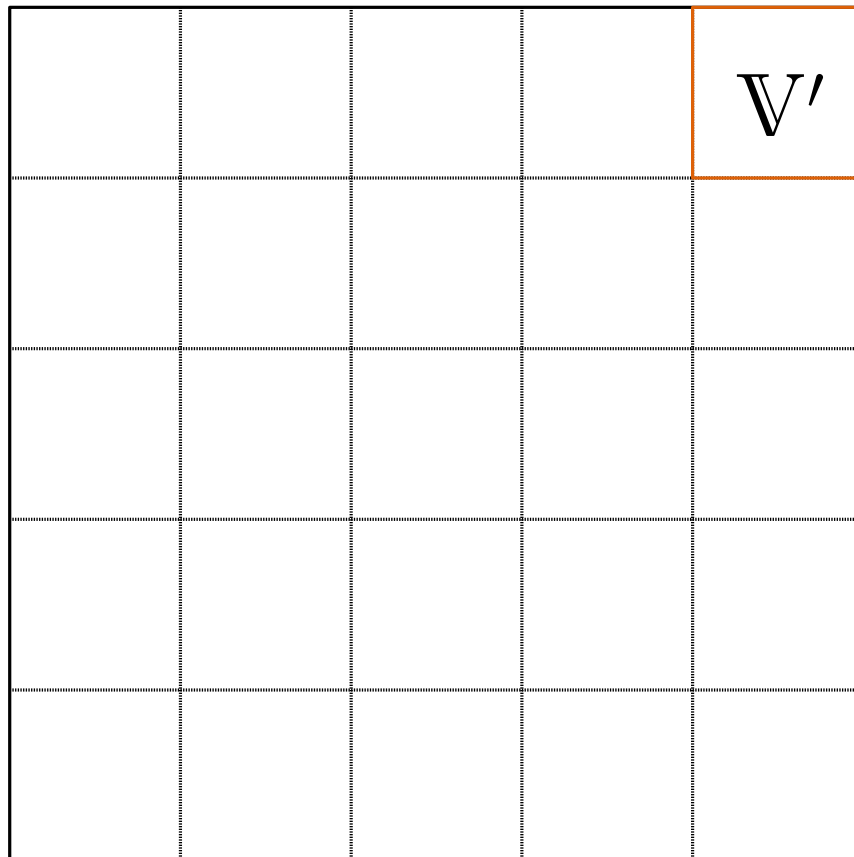
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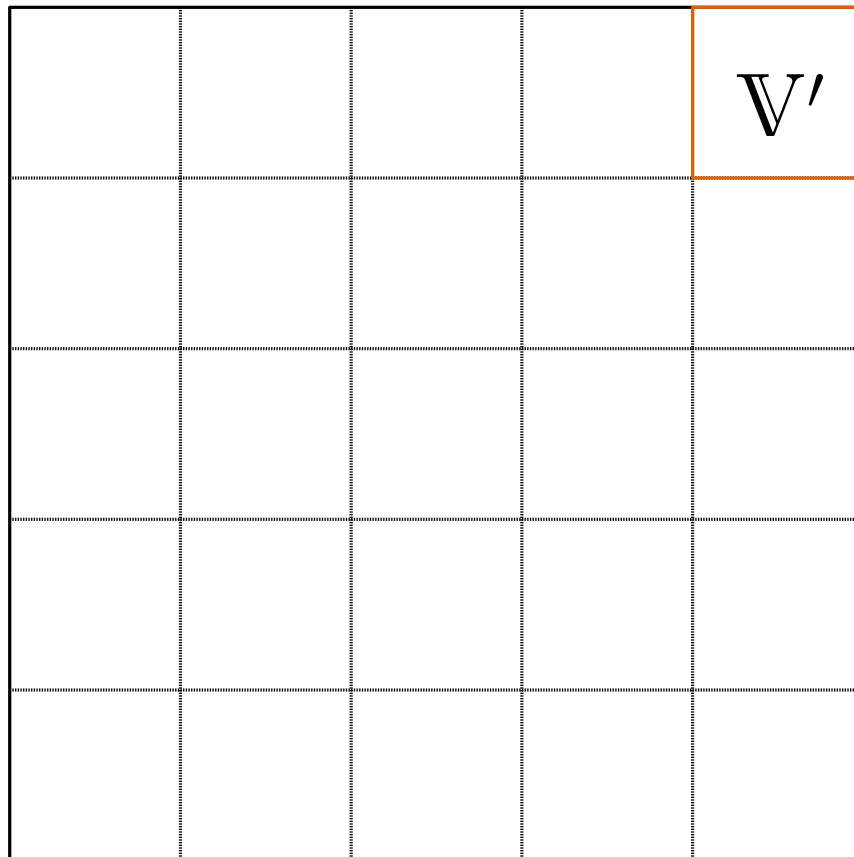


$$\begin{array}{c} \updownarrow \\ < \frac{2r}{\sqrt{d}} \end{array}$$

$\forall x, y \in V'$ it holds that $d(x, y) < 2r$

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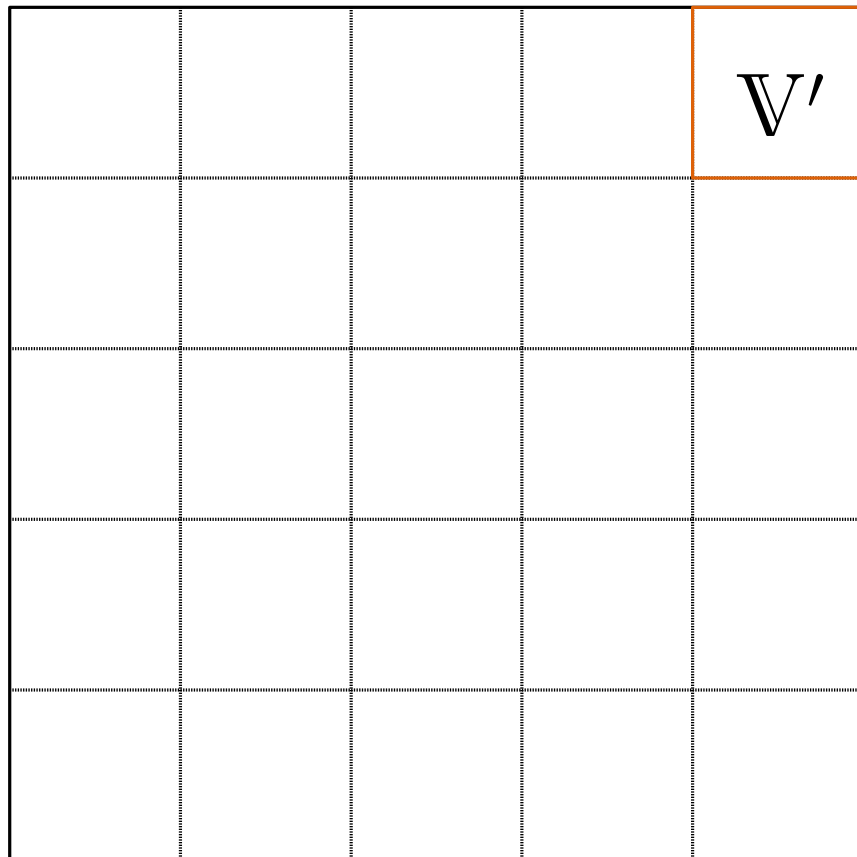
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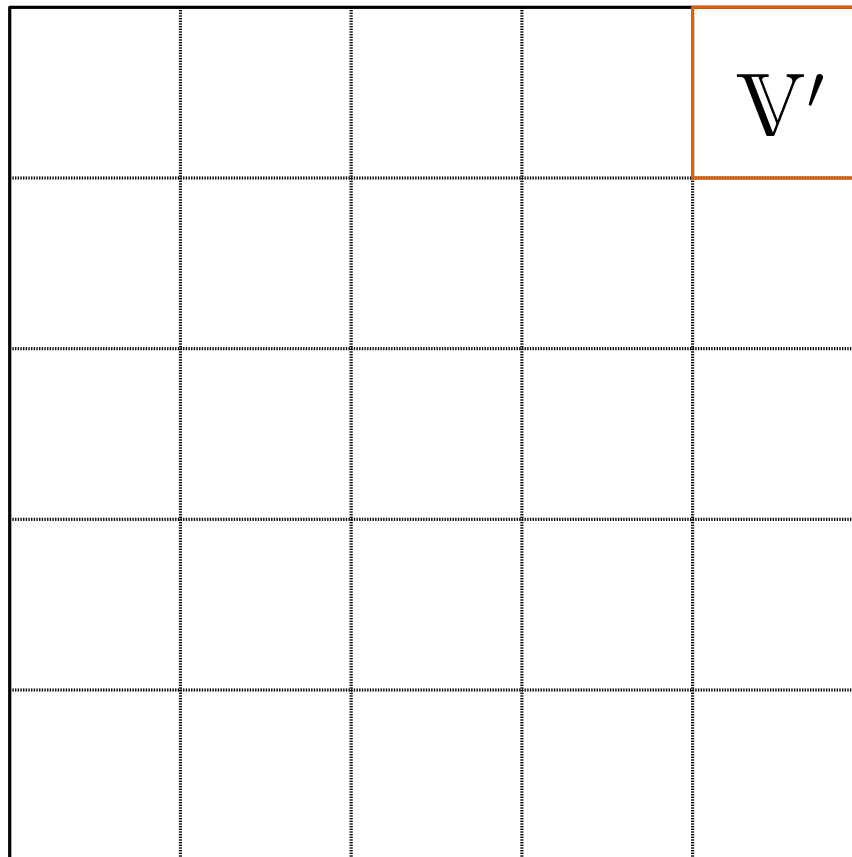
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**These vertices form a clique in G_ρ
for all ρ !**

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\mathbb{V}



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natural clique cover of size $\Theta(\text{vol}(\mathbb{V}))$

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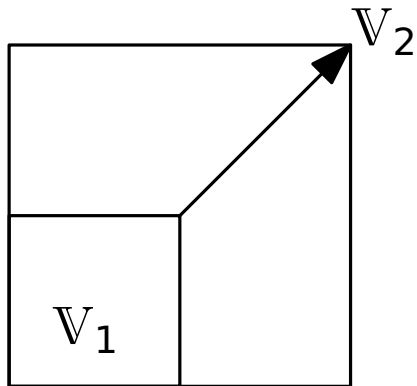
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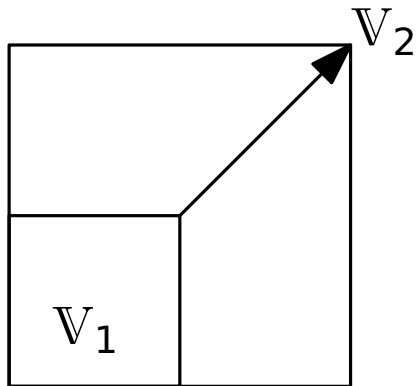
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Theorem: Clique dynamics for a clique cover of size m converge in time $\text{poly}(m)$ for $\lambda < \lambda_c(\Delta)$.

2nd Result: Properties of the partition function

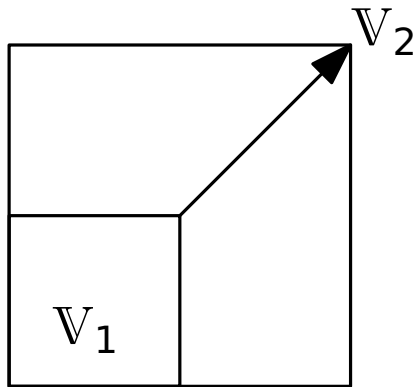


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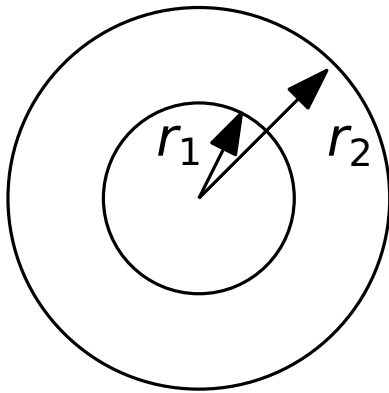


$$\Xi_{V_1}(\lambda, r) \leq \Xi_{V_2}(\lambda, r)$$

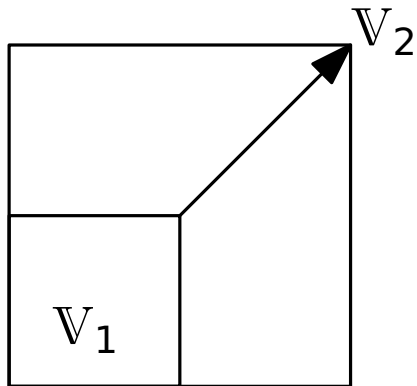
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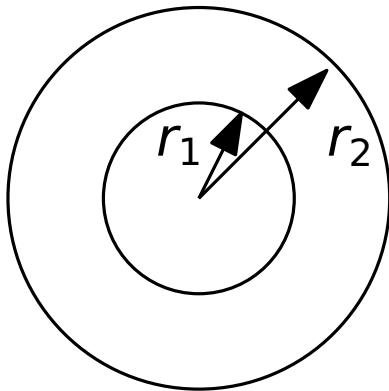
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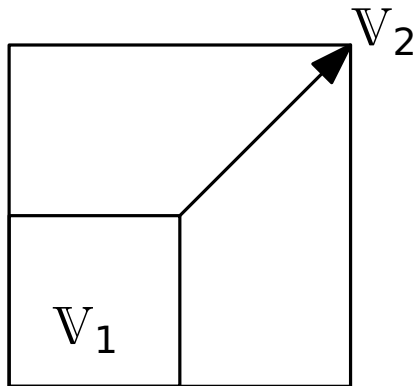


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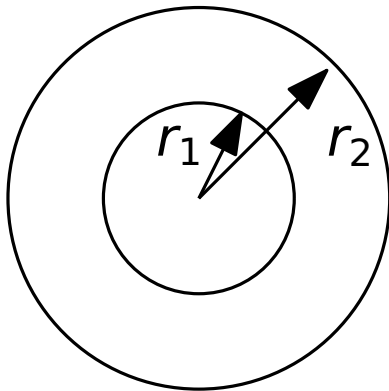


$$\Xi_V(\lambda, r_2) \leq \Xi_V(\lambda, r_1)$$

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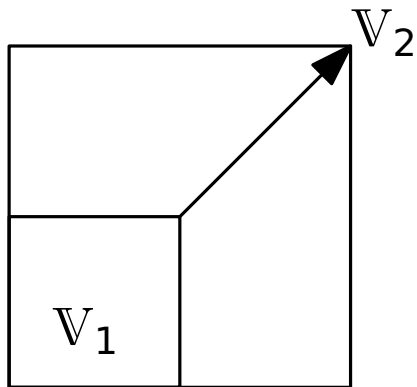
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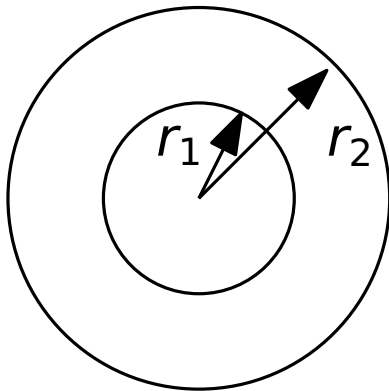
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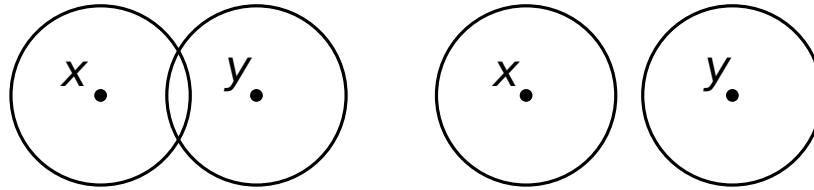


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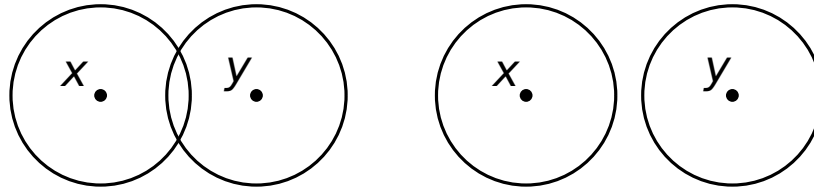
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invalid to valid \longrightarrow increase $r \mapsto r_+$
valid to invalid \longleftarrow decrease $r \mapsto r_-$

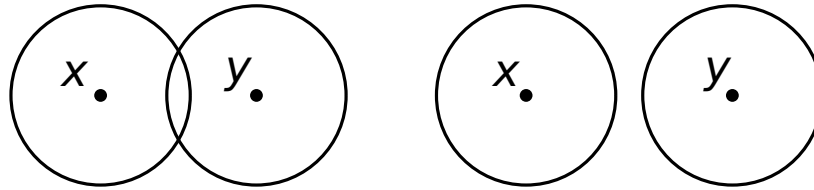
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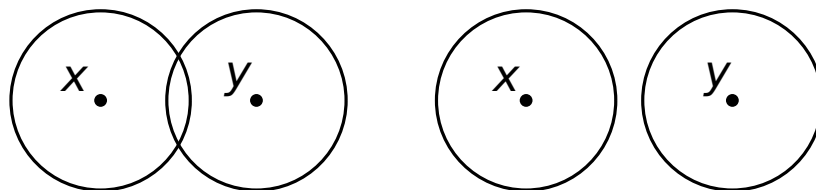
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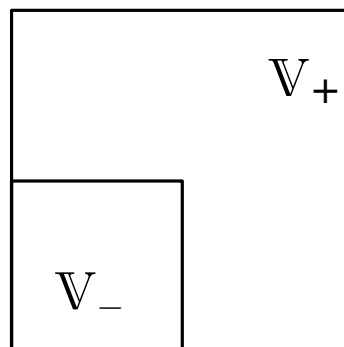


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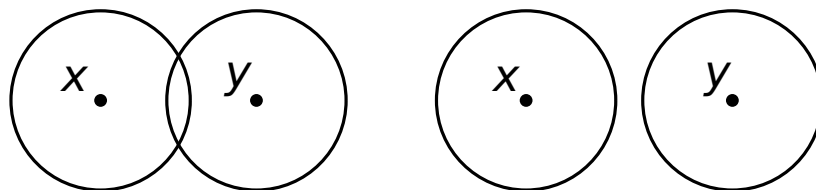
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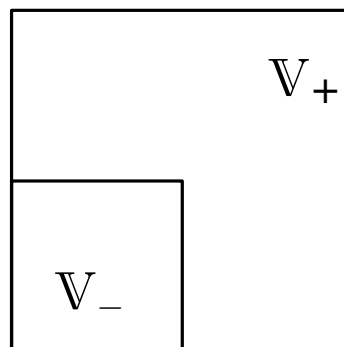


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Theorem: $|\Xi_{\mathbb{V}}(\lambda, r) - Z(G_\rho, \lambda_\rho)| \leq \frac{\Theta(\nu(\mathbb{V})^2)}{\rho} \cdot \Xi_{\mathbb{V}}(\lambda, r)$

2nd Result: Out of the box algorithms

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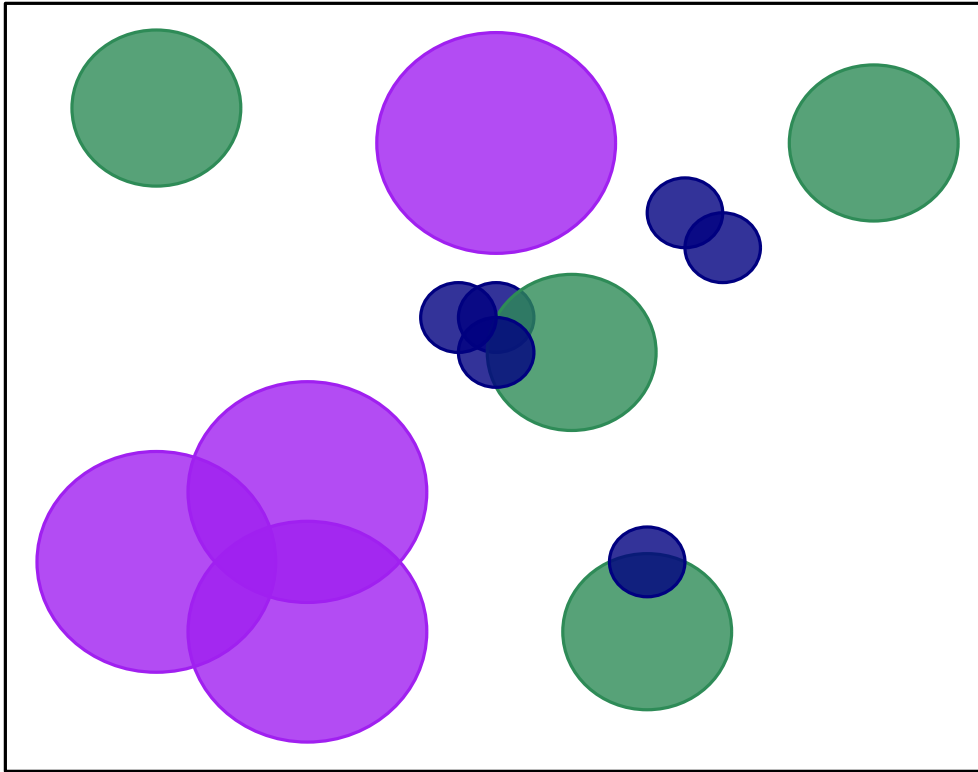
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Approximate sampling:

1. Use an approximate sampler to obtain an independent set of G_{ρ}
2. Recover the points in \mathbb{V} that correspond to the vertices of the independent set
3. Randomly perturb the positions of these points

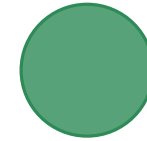
2nd Result: Multiple types of particles



type 1



type 2

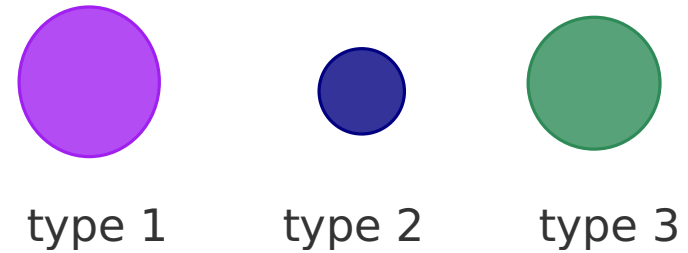
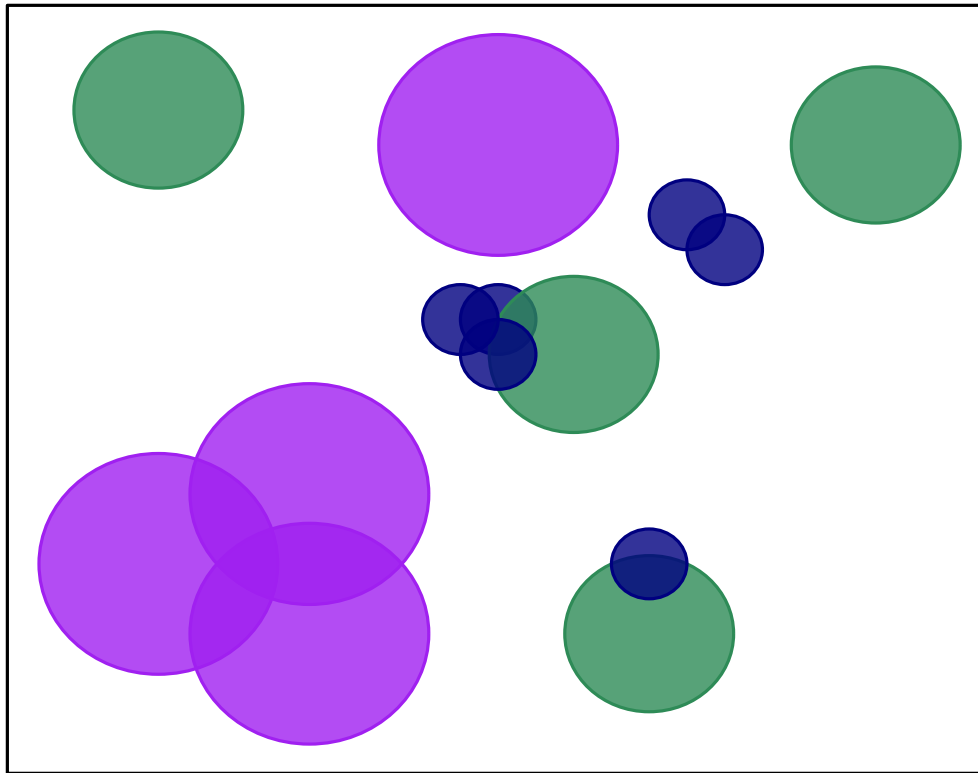


type 3

Interaction Matrix

$$\begin{bmatrix} 0 & r_1 + r_2 & r_1 + r_3 \\ r_1 + r_2 & 0 & 0 \\ r_1 + r_3 & 0 & 2r_3 \end{bmatrix}$$

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Approximation + Sampling Algorithms for :

$$\lambda_{\max} < \frac{e}{\|B\|_1}$$

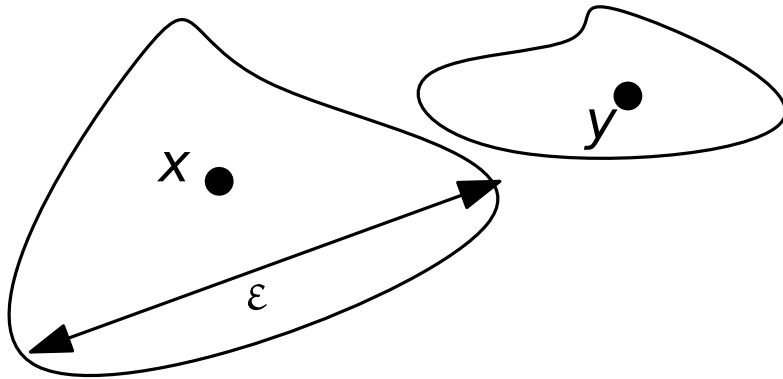
Where B is the volume exclusion matrix, with entries $v(B_{r_{ij}})$

2nd Result: Random point allocations

The previous arguments work for any δ - ε -allocation $\Phi : \mathbb{V} \rightarrow X$

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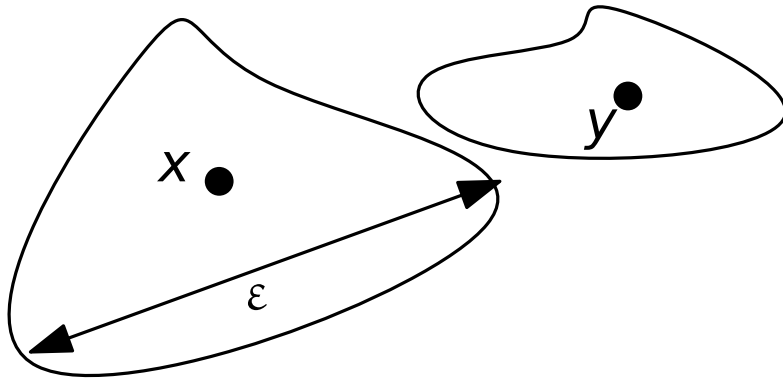
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$$(1 - \delta) \frac{\nu(\mathbb{V})}{|X|} \leq \nu(\Phi^{-1}(x)) \leq (1 + \delta) \frac{\nu(\mathbb{V})}{|X|}$$

2nd Result: Random point allocations

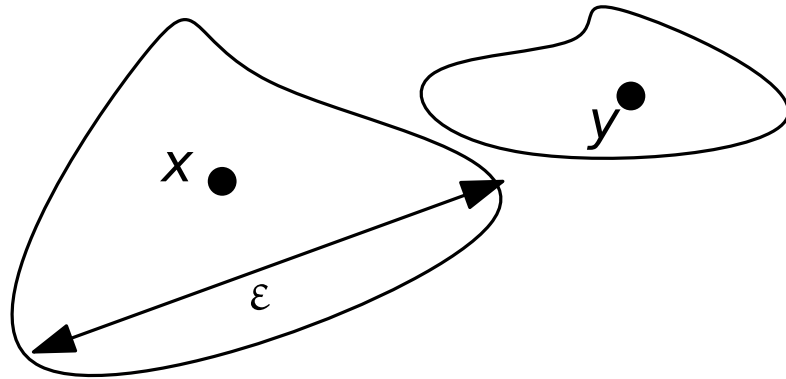
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Theorem:

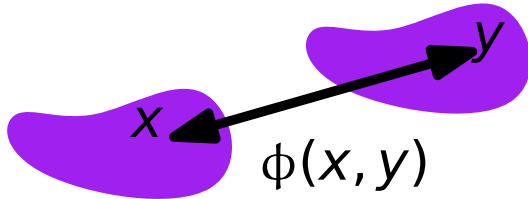
The hard-core model with fugacity $\lambda \nu(\mathbb{V})/n$
of (\mathbb{V}, r) -geometric random graphs

concentrates around $\Xi_{\mathbb{V}}(\lambda, r)$

Spatial process, where the particles interact via repulsive forces

On any complete, separable measure space \mathbb{X} .

$$\phi(x, y) \geq 0$$

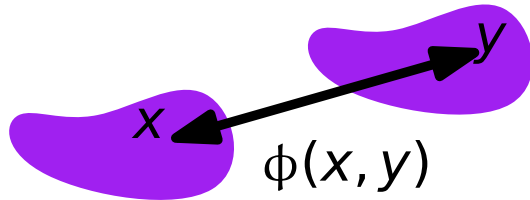


$$H(x_1, \dots, x_k) = \sum_{\{i, j\} \in \binom{[k]}{2}} \phi(x_i, x_j)$$

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Gibbs measure:

$$\frac{dP_{\mathbb{V}}^{(\lambda, \phi)}}{dQ_{\lambda}}(x_1, \dots, x_k) = \frac{\mathbf{1}_{\forall i \in [k]: x_i \in \mathbb{V}} \cdot e^{-H(x_1, \dots, x_k)} e^{\lambda \nu(\mathbb{V})}}{\Xi_{\mathbb{V}}(\lambda, \phi)}$$

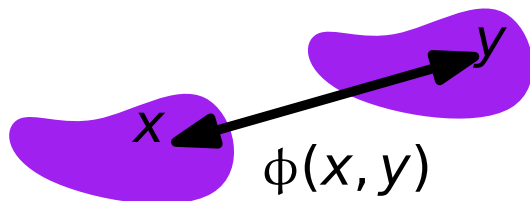
Partition function:

$$\Xi_{\mathbb{V}}(\lambda, \phi) = 1 + \sum_{k \in \mathbb{N}_{\geq 1}} \frac{\lambda^k}{k!} \int_{\mathbb{V}^k} e^{-H(x_1, \dots, x_k)} \nu^k(d\mathbf{x})$$

Spatial process, where the particles interact via repulsive forces

On any complete, separable measure space \mathbb{X} .

$$\phi(x, y) \geq 0$$



$$H(x_1, \dots, x_k) = \sum_{\{i,j\} \in \binom{[k]}{2}} \phi(x_i, x_j)$$

Gibbs measure:

$$\frac{dP_{\mathbb{V}}^{(\lambda, \phi)}}{dQ_{\lambda}}(x_1, \dots, x_k) = \frac{\mathbf{1}_{\forall i \in [k]: x_i \in \mathbb{V}} \cdot e^{-H(x_1, \dots, x_k)} e^{\lambda \nu(\mathbb{V})}}{\Xi_{\mathbb{V}}(\lambda, \phi)}$$

Partition function:

$$\Xi_{\mathbb{V}}(\lambda, \phi) = 1 + \sum_{k \in \mathbb{N}_{\geq 1}} \frac{\lambda^k}{k!} \int_{\mathbb{V}^k} e^{-H(x_1, \dots, x_k)} \nu^k(d\mathbf{x})$$

Hard-sphere model:

$$\phi(x, y) = \begin{cases} \infty, & \text{if } \text{dist}(x, y) < r \\ 0, & \text{otherwise} \end{cases}$$

Temperedness constant: $C_\phi = \text{ess sup}_x \int_{\mathbb{X}} |1 - e^{-\phi(x,y)}| \nu(dy)$

measures the strength of interactions between points

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 - Analyticity of the pressure
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Question: Can we get efficient approximation and sampling algorithms?

The model $\zeta_{\mathbb{V}, \phi}^{(n)}$

- n vertices and bounded measurable region $\mathbb{V} \subseteq \mathbb{X}$
- for each $i \in [n]$ draw a point $x_i \sim \text{unif}_{\mathbb{V}}$ independently
- For all $i, j \in [n]$, with $i \neq j$, connect i and j with an edge with probability $p_{\phi} = 1 - e^{-\phi(x_i, x_j)}$ independently

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Encompasses:

- Erdős–Rényi random graphs
- Geometric random graphs
- Hyperbolic random graphs

Result 3: Hard-core model on $\zeta_{\mathbb{V},\phi}^{(n)}$

Can we show that the hard-core model on $G \sim \zeta_{\mathbb{V},\phi}^{(n)}$ with fugacity $\lambda_{\mathbb{V}}(\mathbb{V})/n$ concentrates around $\Xi_{\mathbb{V}}(\lambda, \phi)$, its expected value?

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Randomized approximation for the partition function when $\lambda < \frac{e}{C_{\phi}}$

Result 3: Approximate sampler

1. Sample the graph G from $\zeta_{\mathbb{V},\phi}^{(n)}$, with $n \in \Theta(v(\mathbb{V}))$
2. For each $v \in V(G)$, keep its position $x_v \in \mathbb{V}$
3. Sample an independent set I from $Z(G, \lambda_{v(\mathbb{V})}/n)$
4. Return the point set X_I that corresponds to the vertices of I .

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To prove that the two densities have small total variation density, we compare each one of them to a Poisson point process of intensity λ utilizing a theorem of Rényi-Mönch

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- Can we get deterministic approximation of Ξ in $\text{poly}(v(\mathbb{V}))$?
- Can we get approximation for $\lambda < \frac{e}{\Delta_\phi}$ without finite range assumption?
- What about other potentials (e.g. Lennard–Jones)?
- Is there a way to show hardness or approximation for some parameter range?